Plug-and-Play Decentralized Model Predictive Control *

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Abstract

In this paper we consider a linear system structured into physically coupled subsystems and propose a decentralized control scheme capable to guarantee asymptotic stability and satisfaction of constraints on system inputs and states. The design procedure is totally decentralized, since the synthesis of a local controller uses only information on a subsystem and its neighbors, i.e. subsystems coupled to it. We first derive tests for checking if a subsystem can be plugged into (or unplugged from) an existing plant without spoiling overall stability and constraint satisfaction. When this is possible, we show how to automatize the design of local controllers so that it can be carried out in parallel by smart actuators equipped with computational resources and capable to exchange information with neighboring subsystems. In particular, local controllers exploit tube-based Model Predictive Control (MPC) in order to guarantee robustness with respect to physical coupling among subsystems. Finally, an application of the proposed control design procedure to frequency control in power networks is presented.

Key Words: Decentralized Control; Decentralized Synthesis; Large-scale Systems; Model Predictive Control; Plug-and-Play Control; Robust Control

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1 Introduction

Decentralized regulators have been studied since the 70's as a viable solution to the control of large-scale systems composed by several physically coupled subsystems [BL88, Lun92]. Compared to centralized schemes, decentralized control structures offer several advantages such as parallel computation of control variables, local transmission of information between each subsystem and the corresponding regulator and higher reliability in presence of controller faults. The problem of guaranteeing stability and suitable performance levels for decentralized control systems has been addressed in the 70's and 80's mainly for unconstrained systems [WD73, Lun92]. Similar remarks apply to distributed control (also known as overlapping decentralized control), where controllers can exchange pieces of information through a communication network (see, e.g. [LA08], and references therein).

Decentralized and distributed control schemes for constrained systems have been proposed only much more recently in the context of Model Predictive Control (MPC) [CJKT02, KBB06, RMS07, FS11, FS12, RFT12, RM09]. These results are particularly appealing because replace large-scale optimization problems stemming from centralized MPC with several smaller-scale problems that can be solved in parallel using computational resources collocated with sensors. While the main focus of decentralized and distributed control is on limiting the computational burden and communication cost associated to real-time cooperations of the control system, attention has also been paid to the complexity of the controller design procedure. In this respect, decentralized and distributed controllers can be designed either in a centralized fashion, i.e. relying on the knowledge of the collective model, or in a decentralized fashion, i.e. not requiring the knowledge of the collective model [BL88, Lun92]. However, decentralized design does not prevent from using collective quantities, based on pieces of information from all subsystems. An example are decentralized control schemes that rely on vector Lyapunov functions for assessing the stability of the closed-loop system [Lun92] and hence require stability analysis of an M -th order system where M is the number of subsystems.

In this paper we move one step further and propose decentralized MPC (DeMPC) schemes with Plug-and-Play (P&P) capabilities. Similarly to [Sto09], P&P means that

- (i) the design of a single controller involves at most information about the subsystem under control and its neighbors, i.e. no step of the design procedure involves collective quantities;
- (ii) when a subsystem joins/leaves an existing plant there is a procedure for
  - (a) assessing if the operation does not spoil stability and constraint satisfaction for the overall plant;
  - (b) automatically retuning at most the controllers of the subsystem and its successors, i.e. subsystems influenced by it.

P&P controllers are attractive for the following reasons. First, the complexity of designing a controller for a given subsystem scales with the number of neighboring subsystem only. Second, P&P eases the revamping of control systems by enabling the replacement of actuators with limited interaction of human operators. It is well known that, for general interconnection topologies, requirement (i) above implies the design of regulators for each subsystem that are robust to the coupling with neighboring subsystems [Lun92]. Our design procedure is no exception and we will exploit tube-based MPC [MSR05] for the design of robust local controllers. While this introduces

an unavoidable degree of conservatism, we argue that P&P DeMPC can be successfully applied in a number of real world plants where coupling among subsystems is sufficiently weak. As an example, we will use P&P DeMPC for designing the Automatic Generation Control (AGC) layer for frequency control in a realistic power network and discuss the plugging in and unplugging of generators areas.

The paper is structured as follows. The design of decentralized controllers is introduced in Section 2 with a focus on the assumptions needed for guaranteeing asymptotic stability of the origin and constraint satisfaction. In Section 3 we discuss how to design the local controllers in a distributed fashion and in Section 4 we describe P&P operations. In Section 5 we discuss the practical design of the local controllers. In Section 6 we present the application of P&P DeMPC to frequency control in a power network and Section 7 is devoted to concluding remarks.

Notation. We use $a : b$ for the set of integers $\{a, a + 1, \dots, b\}$. The column vector with s components $v_1, \dots, v_s$ is $\mathbf{v} = (v_1, \dots, v_s)$. The function $\text{diag}(G_1, \dots, G_s)$ denotes the block-diagonal matrix composed by s block G_i , $i \in 1 : s$. The pseudo-inverse of a matrix $A \in \mathbb{R}^{m \times n}$ is denoted with $A^\dagger$. The symbol $\oplus$ denotes the Minkowski sum of sets, i.e. $A = B \oplus C$ if and only if $A = \{a : a = b + c, \text{ for all } b \in B \text{ and } c \in C\}$. Moreover, $\bigoplus_{i=1}^s G_i = G_1 \oplus \dots \oplus G_s$. The symbols $\mathbf{1}_n$ and $\mathbf{0}_n$ denote the column vectors with n elements equal to 1 and 0, respectively. A zonotope is a centrally symmetric convex polytopes: given a vector $p \in \mathbb{R}^n$ and a matrix $\Xi \in \mathbb{R}^{n \times m}$, the zonotope $\mathbb{X} \subseteq \mathbb{R}^n$ is the set $\mathbb{X} = \{x \mid x = p + \Xi d, \|d\|_\infty \leq 1\}$, with $d \in \mathbb{R}^m$. Moreover, if $\mathbb{X}$ is a zonotope, its support function in the direction $c \in \mathbb{R}^n$ is given by [KG98] as

$$\sup_{x \in \mathbb{X}} c^T x = \| \Xi^T c \|_1. \quad (1)$$

The set $\mathbb{X} \subseteq \mathbb{R}^n$ is Robust Positively Invariant (RPI) [RM09] for $x(t+1) = f(x(t), w(t))$, $w(t) \in \mathbb{W} \subseteq \mathbb{R}^m$ if $x(t) \in \mathbb{X} \Rightarrow f(x(t), w(t)) \in \mathbb{X}$, $\forall w(t) \in \mathbb{W}$. The RPI set $\underline{\mathbb{X}}$ is minimal if every other RPI $\mathbb{X}$ verifies $\underline{\mathbb{X}} \subseteq \mathbb{X}$. The RPI set $\mathbb{X}(\delta)$ is a δ -outer approximation of the minimal RPI $\underline{\mathbb{X}}$ if

$$x \in \mathbb{X}(\delta) \Rightarrow \exists \underline{x} \in \underline{\mathbb{X}} \text{ and } \sigma \in B_\delta(0) : x = \underline{x} + \sigma. \quad (2)$$

where $B_\delta(v)$ is the 2-norm open ball of radius δ centered in $v \in \mathbb{R}^n$.

2 Decentralized tube-based MPC of linear systems

We consider a discrete-time Linear Time-Invariant (LTI) system

$$\mathbf{x}^+ = \mathbf{A}\mathbf{x} + \mathbf{B}\mathbf{u} \quad (3)$$

where $\mathbf{x} \in \mathbb{R}^n$ and $\mathbf{u} \in \mathbb{R}^m$ are the state and the input, respectively, at time t and $\mathbf{x}^+$ stands for $\mathbf{x}$ at time $t+1$. We will use the notation $\mathbf{x}(t)$, $\mathbf{u}(t)$ only when necessary. The state is partitioned into M state vectors $x_{[i]} \in \mathbb{R}^{n_i}$, $i \in \mathcal{M} = 1 : M$ such that $\mathbf{x} = (x_{[1]}, \dots, x_{[M]})$, and $n = \sum_{i \in \mathcal{M}} n_i$. Similarly, the input is partitioned into M vectors $u_{[i]} \in \mathbb{R}^{m_i}$, $i \in \mathcal{M}$ such that $\mathbf{u} = (u_{[1]}, \dots, u_{[M]})$ and $m = \sum_{i \in \mathcal{M}} m_i$.

We assume the dynamics of the i -th subsystem is given by

$$\Sigma_{[i]} : \quad x^+_{[i]} = A_{ii}x_{[i]} + B_i u_{[i]} + w_{[i]} \quad (4)$$

$$w_{[i]} = \sum_{j \in \mathcal{N}_i} A_{ij}x_{[j]} \quad (5)$$

where $A_{ij} \in \mathbb{R}^{n_i \times n_j}$, $i, j \in \mathcal{M}$, $B_i \in \mathbb{R}^{n_i \times m_i}$ and $\mathcal{N}_i$ is the set of neighbors to subsystem i defined as $\mathcal{N}_i = \{j \in \mathcal{M} : A_{ij} \neq 0, i \neq j\}$.

According to (4), the matrix $\mathbf{A}$ in (3) is decomposed into blocks A_{ij} , $i, j \in \mathcal{M}$. We also define $\mathbf{A}_D = \text{diag}(A_{11}, \dots, A_{MM})$ and $\mathbf{A}_C = \mathbf{A} - \mathbf{A}_D$, i.e. $\mathbf{A}_D$ collects the state transition matrices of every subsystem and $\mathbf{A}_C$ collects coupling terms between subsystems. From (4) one also obtains $\mathbf{B} = \text{diag}(B_1, \dots, B_M)$ because submodels (4) are input decoupled.

In this Section we propose a decentralized controller for (3) guaranteeing asymptotic stability of the origin of the closed-loop system and constraints satisfaction.

In the spirit of tube-based control [MSR05], we treat $w_{[i]}$ as a disturbance and equip (4) with the controller $\mathcal{C}_{[i]}$ given by

$$u_{[i]} = v_{[i]} + K_i(x_{[i]} - \bar{x}_{[i]}). \quad (6)$$

where $K_i \in \mathbb{R}^{m_i \times n_i}$, $i \in \mathcal{M}$ and variables $v_{[i]}$ and $\bar{x}_{[i]}$ will be computed by a local state-feedback MPC controller, i.e. there exist functions $\kappa_i : \mathbb{R}^{n_i} \rightarrow \mathbb{R}^{m_i}$ and $\eta_i : \mathbb{R}^{n_i} \rightarrow \mathbb{R}^{n_i}$ such that $v_{[i]} = \kappa_i(x_{[i]})$ and $\bar{x}_{[i]} = \eta_i(x_{[i]})$. Note that the controller $\mathcal{C}_{[i]}$ is completely decentralized, since it depends upon quantities of system $\Sigma_{[i]}$ only.

Next, we clarify properties of matrices K_i , $i \in \mathcal{M}$ that are required for the stability of system (3) controlled by (6). Defining the collective variables $\bar{\mathbf{x}} = (\bar{x}_{[1]}, \dots, \bar{x}_{[M]}) \in \mathbb{R}^n$, $\mathbf{v} = (v_{[1]}, \dots, v_{[M]}) \in \mathbb{R}^m$ and the matrix $\mathbf{K} = \text{diag}(K_1, \dots, K_M) \in \mathbb{R}^{m \times n}$, from (4) and (6) one obtains the collective model

$$\mathbf{x}^+ = (\mathbf{A} + \mathbf{BK})\mathbf{x} + \mathbf{B}(\mathbf{v} - \mathbf{K}\bar{\mathbf{x}}). \quad (7)$$

The following assumptions will be needed for designing stabilizing controllers $\mathcal{C}_{[i]}$.

Assumption 1. (i) The matrices $F_i = A_{ii} + B_i K_i$, $i \in \mathcal{M}$ are Schur.

(ii) The matrix $\mathbf{F} = \mathbf{A} + \mathbf{BK}$ is Schur.

■

We discuss now constraints satisfaction. To this purpose, we equip subsystems $\Sigma_{[i]}$, $i \in \mathcal{M}$ with the constraints $x_{[i]} \in \mathbb{X}_i$, $u_{[i]} \in \mathbb{U}_i$, define the sets $\mathbb{X} = \prod_{i \in \mathcal{M}} \mathbb{X}_i$, $\mathbb{U} = \prod_{i \in \mathcal{M}} \mathbb{U}_i$ and consider the collective constrained system (3) with

$$\mathbf{x} \in \mathbb{X}, \mathbf{u} \in \mathbb{U}. \quad (8)$$

As in tube-based MPC control [MSR05], our goal is to compute tightened state constraints $\hat{\mathbb{X}}_i \subseteq \mathbb{X}_i$ and input constraints $\mathbb{V}_i \subseteq \mathbb{U}_i$ guaranteeing that

$$\begin{aligned} \bar{x}_{[i]}(k) &\in \hat{\mathbb{X}}_i, v_{[i]}(k) \in \mathbb{V}_i, \forall i \in \mathcal{M} \\ &\Rightarrow \mathbf{x}(k+1) \in \mathbb{X}, \mathbf{u}(k) \in \mathbb{U}, \end{aligned} \quad (9)$$

The next Assumption characterizes the shape of constraints $\mathbb{X}_i$, $\hat{\mathbb{X}}_i$, $\mathbb{U}_i$ and $\mathbb{V}_i$, $i \in \mathcal{M}$.

Assumption 2. Constraints $\mathbb{X}_i$ and $\hat{\mathbb{X}}_i$ are zonotopes given by

$$\begin{aligned} \mathbb{X}_i &= \{x_{[i]} \in \mathbb{R}^{n_i} | f_{i,r}^T x_{[i]} \leq 1, r \in 1 : \bar{r}_i\} = \{x_{[i]} \in \mathbb{R}^{n_i} | \mathcal{F}_i x_{[i]} \leq \mathbf{1}_{\bar{r}_i}\} \\ &= \{x_{[i]} \in \mathbb{R}^{n_i} | x_{[i]} = \Xi_i d_i, \|d_i\|_\infty \leq 1\}, \end{aligned} \quad (10)$$

$$\begin{aligned} \hat{\mathbb{X}}_i &= \{\hat{x}_{[i]} \in \mathbb{R}^{n_i} | \hat{f}_{i,r}^T \hat{x}_{[i]} \leq \hat{l}_i, r \in 1 : \bar{r}_i\} = \{\hat{x}_{[i]} \in \mathbb{R}^{n_i} | \hat{\mathcal{F}}_i \hat{x}_{[i]} \leq \hat{l}_i \mathbf{1}_{\bar{r}_i}\} \\ &= \{\hat{x}_{[i]} = \hat{\Xi}_i \hat{d}_i, \|\hat{d}_i\|_\infty \leq \hat{l}_i\}, \end{aligned} \quad (11)$$

where $\mathcal{F}_i = (f_{i,1}^T, \dots, f_{i,\bar{r}_i}^T) \in \mathbb{R}^{\bar{r}_i \times n_i}$, $\text{rank}(\mathcal{F}_i) = n_i$, $d_i \in \mathbb{R}^{e_i}$, $\Xi_i \in \mathbb{R}^{n_i \times e_i}$, $\hat{l}_i \in \mathbb{R}_+$, $\hat{\mathcal{F}}_i = (\hat{f}_{i,1}^T, \dots, \hat{f}_{i,\hat{r}_i}^T) \in \mathbb{R}^{\hat{r}_i \times n_i}$, $\hat{d}_i \in \mathbb{R}^{\bar{e}_i}$ and $\hat{\Xi}_i \in \mathbb{R}^{n_i \times \bar{e}_i}$.

Constraints $\mathbb{U}_i$ and $\mathbb{V}_i$, $i \in \mathcal{M}$ are polytopes containing the origin in their interior, that, without loss of generality, are defined as follows

$$\mathbb{U}_i = \{u_{[i]} \in \mathbb{R}^{m_i} | h_{i,r}^T u_{[i]} \leq 1, r \in 1 : r_{u_i}\} = \{u_{[i]} \in \mathbb{R}^{m_i} | \mathcal{H}_i u_{[i]} \leq \mathbf{1}_{r_{u_i}}\}, \quad (12)$$

$$\mathbb{V}_i = \{v_{[i]} \in \mathbb{R}^{m_i} | h_{i,r}^T v_{[i]} \leq 1 - l_{v_i,r}, r \in 1 : r_{u_i}\} = \{v_{[i]} \in \mathbb{R}^{m_i} | \mathcal{H}_i v_{[i]} \leq \mathbf{1}_{r_{u_i}} - l_{v_i}\}, \quad (13)$$

where $\mathcal{H}_i = (h_{i,1}^T, \dots, h_{i,r_{u_i}}^T) \in \mathbb{R}^{r_{u_i} \times m_i}$, $l_{v_i,r} \in \mathbb{R}_+$ and $l_{v_i} = (l_{v_i,1}, \dots, l_{v_i,r_{u_i}})$.

■

From the results in [KG98], under Assumptions 1-(i) and 2 there exist nonempty RPIs $\mathbb{Z}_i \subseteq \mathbb{R}^{n_i}$, $i \in \mathcal{M}$ for the dynamics

$$z^+_{[i]} = (A_{ii} + B_i K_i) z_{[i]} + w_{[i]} \quad (14)$$

and $w_i \in \mathbb{W}_i = \bigoplus_{j \in \mathcal{N}_i} A_{ij} \mathbb{X}_j$. In particular, for $\delta_i > 0$, we denote with $\mathbb{Z}_i(\delta_i)$ an RPI set that is a δ_i -outer approximation of the minimal RPI for (14) and $w_{[i]} \in \mathbb{W}_i$.

For guaranteeing (9), we introduce the following Assumption.

Assumption 3. *There exist $\delta_i > 0$ and nonempty constraint sets $\hat{\mathbb{X}}_i$ and $\mathbb{V}_i$, $\forall i \in \mathcal{M}$ verifying*

$$\hat{\mathbb{X}}_i \oplus \mathbb{Z}_i(\delta_i) \subseteq \mathbb{X}_i \quad (15)$$

$$\mathbb{V}_i \oplus K_i \mathbb{Z}_i(\delta_i) \subseteq \mathbb{U}_i. \quad (16)$$

■

Note that, by construction, one has $\mathbb{Z}_i(\delta_i) \supseteq \mathbb{W}_i$ and therefore (15) and (16) cannot be verified if $\mathbb{W}_i$ is “too big”, i.e. $\mathbb{W}_i \supseteq \mathbb{X}_i$ or $K_i \mathbb{W}_i \supseteq \mathbb{U}_i$.

Under Assumptions 1-3, as in [MSR05], we set in (6)

$$\kappa_i(x_{[i]}(t)) = v_{[i]}(0|t), \quad \eta_i(x_{[i]}(t)) = \hat{x}_{[i]}(0|t) \quad (17)$$

where $v_{[i]}(0|t)$ and $\hat{x}_{[i]}(0|t)$ are optimal values of variables $v_{[i]}(0)$ and $\hat{x}_{[i]}(0)$, respectively, appearing in the following MPC- i problem to be solved at time t

$$\mathbb{P}_i^N(x_{[i]}(t)) = \min_{\substack{\hat{x}_{[i]}(0) \\ v_{[i]}(0:N_i-1)}} \sum_{k=0}^{N_i-1} \ell_i(\hat{x}_{[i]}(k), v_{[i]}(k)) + V_{f_i}(\hat{x}_{[i]}(N_i)) \quad (18a)$$

$$x_{[i]}(t) - \hat{x}_{[i]}(0) \in \mathbb{Z}_i(\delta_i) \quad (18b)$$

$$\hat{x}_{[i]}(k+1) = A_{ii} \hat{x}_{[i]}(k) + B_i v_{[i]}(k) \quad k \in 0 : N_i - 1 \quad (18c)$$

$$\hat{x}_{[i]}(k) \in \hat{\mathbb{X}}_i \quad k \in 0 : N_i - 1 \quad (18d)$$

$$v_{[i]}(k) \in \mathbb{V}_i \quad k \in 0 : N_i - 1 \quad (18e)$$

$$\hat{x}_{[i]}(N_i) \in \hat{\mathbb{X}}_{f_i} \quad (18f)$$

In (18), $N_i \in \mathbb{N}$ is the prediction horizon, $\ell_i(\hat{x}_{[i]}(k), v_{[i]}(k)) : \mathbb{R}^{n_i \times m_i} \rightarrow \mathbb{R}_+$ is the stage cost and $V_{f_i}(\hat{x}_{[i]}(N_i)) : \mathbb{R}^{n_i} \rightarrow \mathbb{R}_+$ is the final cost, fulfilling the following assumption.

Assumption 4. *For all $i \in \mathcal{M}$, there exist an auxiliary control law $\kappa_i^{aux}(\hat{x}_{[i]})$ and a $\mathcal{K}_\infty$ function $\mathcal{B}_i$ such that:*

- (i) $\ell_i(x_{[i]}, u_{[i]}) \geq \mathcal{B}_i(\|(x_{[i]}, u_{[i]})\|)$, for all $x_{[i]} \in \mathbb{R}^{n_i}$, $u_{[i]} \in \mathbb{R}^{m_i}$ and $\ell_i(0, 0) = 0$;
- (ii) $\hat{\mathbb{X}}_{f_i} \subseteq \hat{\mathbb{X}}_i$ is an invariant set for $\hat{x}_{[i]}^+ = A_{ii}\hat{x}_{[i]} + B_i\kappa_i^{aux}(\hat{x}_{[i]})$;
- (iii) $\forall \hat{x}_{[i]} \in \hat{\mathbb{X}}_{f_i}$, $\kappa_i^{aux}(\hat{x}_{[i]}) \in \mathbb{V}_i$;
- (iv) $\forall \hat{x}_{[i]} \in \hat{\mathbb{X}}_{f_i}$, $V_{f_i}(\hat{x}_{[i]}^+) - V_{f_i}(\hat{x}_{[i]}) \leq -\ell_i(\hat{x}_{[i]}, \kappa_i^{aux}(\hat{x}_{[i]}))$.

■

We highlight that there are several methods, discussed e.g. in [RM09], for computing $\ell_i(\cdot)$, $V_{f_i}(\cdot)$ and $\mathbb{X}_{f_i}$ verifying Assumption 4.

The next Theorem, that is proved in Appendix A, provides the main results on stability of the closed-loop system (7) and (17) equipped with constraints (8).

Theorem 1. *Let Assumptions 1-4 hold. Define the feasibility region for the MPC-i problem as*

$$\mathbb{X}_i^N = \{s_{[i]} \in \mathbb{X}_i : (18) \text{ is feasible for } x_{[i]}(t) = s_{[i]}\}$$

and the collective feasibility region as $\mathbb{X}^N = \prod_{i \in \mathcal{M}} \mathbb{X}_i^N$.

Then

- (i) *if $\mathbf{x}(0) \in \mathbb{X}^N$, i.e. $x_{[i]}(0) \in \mathbb{X}_i^N$ for all $i \in \mathcal{M}$, constraints (8) are fulfilled at all time instants;*
- (ii) *the origin of the closed-loop system (7) and (17) is asymptotically stable and $\mathbb{X}^N$ is a region of attraction.*

■

In order to design a DeMPC scheme based on MPC-i problems (18) and for which Theorem 1 applies, the main problem that still has to be solved is the following one.

Problem $\mathcal{P}$

Compute matrices K_i , $i \in \mathcal{M}_i$, if they exist, verifying Assumptions 1 and 3.

■

In the next Section, we show how to solve Problem $\mathcal{P}$ in a distributed fashion under Assumption 2 complemented by the next assumption.

Assumption 5. *Matrices $\hat{\mathcal{F}}_i$ (and hence $\hat{\Xi}_i$) in (11) are given for $i \in \mathcal{M}$.*

■

Note that Assumption 5 fixes the shape of set $\hat{\mathbb{X}}_i$, $i \in \mathcal{M}$ leaving the freedom to choose the zooming parameters $\hat{l}_i$. Also the shape of each set $\mathbb{V}_i$ is fixed and, from Assumption 2, it coincides with the shape of $\mathbb{U}_i$ while parameters l_{v_i} are free.

3 Decentralized synthesis of DeMPC

The next Theorem will allow us to solve Problem $\mathcal{P}$ in a distributed fashion.

Theorem 2. *Let Assumptions 2 and 5 hold. For given matrices K_i , $i \in \mathcal{M}$, verifying Assumption 1-(i), if the following conditions are fulfilled*

$$\alpha_i = \sum_{j \in \mathcal{N}_i} \sum_{k=0}^{\infty} \|\mathcal{F}_i F_i^k A_{ij} \mathcal{F}_j^b\|_{\infty} < 1, \quad \forall i \in \mathcal{M} \quad (19)$$

then

(i) Assumption 1-(ii) holds.

(ii) For all $i \in \mathcal{M}$, defining

$$\hat{\underline{L}}_{i,r} = \frac{1 - \sum_{j \in \mathcal{N}_i} \sum_{k=0}^{\infty} \|f_{i,r}^T F_i^k A_{ij} \Xi_j\|_{\infty}}{\|f_{i,r}^T \hat{\Xi}_i\|_{\infty}}, \quad r \in 1 : \bar{r}_i \quad (20)$$

there is $\delta_i > 0$ such that

$$\hat{L}_i = \min_{r \in 1 : \bar{r}_i} \hat{\underline{L}}_{i,r} - \frac{\|f_{i,r}^T\|_{\infty} \delta_i}{\|f_{i,r}^T \hat{\Xi}_i\|_{\infty}} > 0. \quad (21)$$

Furthermore, choosing $\hat{l}_i \in (0, \hat{L}_i]$ and the set $\hat{\mathbb{X}}_i$ as in (11), the inclusion (15) holds.

(iii) For $\delta_i > 0$ verifying (21) assume the following condition is fulfilled

$$\beta_i(\delta_i) = \max_{r \in 1 : r_{u_i}} \hat{l}_{v_i,r}(\delta_i) < 1 \quad (22)$$

with

$$\hat{l}_{v_i,r}(\delta_i) = \sup_{z_i \in \mathbb{Z}_i(\delta_i)} h_{i,r}^T K_i z_i, \quad r \in 1 : r_{u_i}. \quad (23)$$

Then, choosing $\mathbb{V}_i$ as in (13) for $l_{v_i,r} = \hat{l}_{v_i,r}(\delta_i)$ the inclusion (16) holds. ■

The proof of Theorem 2 can be found in Appendix B.

We highlight that under Assumption 5, for a given $i \in \mathcal{M}$, the quantities α_i in (19), $\hat{L}_i$ in (21) and $\beta_i(\delta_i)$ in (22) depend only upon local fixed parameters $\{A_{ii}, B_i, \mathcal{F}_i, \mathcal{H}_i\}$, neighbors' fixed parameters $\{A_{ij}, \Xi_j\}_{j \in \mathcal{N}_i}$ (or equivalently $\{A_{ij}, \mathcal{F}_j\}_{j \in \mathcal{N}_i}$) and local tunable parameters $\{K_i, \delta_i\}$ but not on neighbors' tunable parameters. Moreover, also the computation of sets $\mathbb{Z}_i(\delta_i)$ depends upon the same parameters. This implies that the choice of $\{K_i, \delta_i\}$ does not influence the choice of $\{K_j, \delta_j\}_{j \neq i}$ and therefore Problem $\mathcal{P}$ is decomposed in the following independent problems for $i \in \mathcal{M}$.

Problem $\mathcal{P}_i$

Check if there exist K_i and $\delta_i > 0$ such that $\alpha_i < 1$, $\hat{L}_i > 0$ and $\beta_i(\delta_i) < 1$. ■

Algorithm 1 Design of controller $\mathcal{C}_{[i]}$ for system $\Sigma_{[i]}$

Input: $A_{ii}, B_i, \mathbb{X}_i, \mathbb{U}_i, \mathcal{N}_i, \{A_{ij}\}_{j \in \mathcal{N}_i}, \{\mathbb{X}_j\}_{j \in \mathcal{N}_i}$

Output: controller $\mathcal{C}_{[i]}$ in (6)

- 1) Find K_i and $\delta_i > 0$ such that Assumption 1-(i) is fulfilled, $\alpha_i < 1$, (21) holds and $\beta_i(\delta_i) < 1$. If they do not exist **stop** (the controller $\mathcal{C}_{[i]}$ cannot be designed).
  - 2) Compute sets $\mathbb{W}_i = \bigoplus_{j \in \mathcal{N}_i} A_{ij} \mathbb{X}_j$ and $\mathbb{Z}_i(\delta_i)$.
  - 3) Compute $\hat{L}_i$ as in (21), choose $\hat{l}_i = \hat{L}_i$ and define $\hat{\mathbb{X}}_i$ as in (11).
  - 4) Compute $\hat{l}_{v_{i,r}}(\delta_i)$ as in (23), set $l_{v_{i,r}} = \hat{l}_{v_{i,r}}(\delta_i)$ and define $\mathbb{V}_i$ as in (13).
  - 5) Compute $\ell_i(\cdot)$, $V_{f_i}(\cdot)$ and $\mathbb{X}_{f_i}$ verifying Assumption 4.
-

According to Theorem 2, the solution to Problem $\mathcal{P}_i$ enables the computation of sets $\hat{\mathbb{X}}_i$ and $\mathbb{V}_i$ and therefore the decentralized design of controller MPC- i . The overall procedure for the decentralized synthesis of local controllers $\mathcal{C}_{[i]}$, $i \in \mathcal{M}$ is summarized in Algorithm 1, whose computational aspects are discussed in Section 5.

In view of the previous discussion, the link between controllers designed through Algorithm 1 and Theorem 1 can be summarized as follows.

Proposition 1. *Under Assumptions 2 and 5, if for all $i \in \mathcal{M}$ controllers $\mathcal{C}_{[i]}$ are designed according to Algorithm 1, then all Assumptions of Theorem 1 are fulfilled.*

4 Plug-and-play operations

In this Section we discuss the synthesis of new controllers and the redesign of existing ones when subsystems are added to or removed from system (4). The goal will be to preserve stability of the origin and constraint satisfaction for the new closed-loop system. Note that plugging in and unplugging of subsystems are here considered as off-line operations. Therefore, the overall plant is not modeled as a switching system. As a starting point, we consider a plant composed by subsystems $\Sigma_{[i]}$, $i \in \mathcal{M}$ equipped with local controllers $\mathcal{C}_{[i]}$, $i \in \mathcal{M}$ produced by Algorithm 1.

4.1 Plugging in operation

We start considering the plugging of subsystem $\Sigma_{[M+1]}$, characterized by parameters $A_{M+1}, B_{M+1}, \mathbb{X}_{M+1}, \mathbb{U}_{M+1}, \mathcal{N}_{M+1}$ and $\{A_{ij}\}_{j \in \mathcal{N}_{M+1}}$, into an existing plant. In particular $\mathcal{N}_{M+1}$ identifies the subsystems that will be physically coupled to $\Sigma_{[M+1]}$ and $\{A_{ij}\}_{j \in \mathcal{N}_{M+1}}$ are the corresponding coupling terms. For building the controller $\mathcal{C}_{[M+1]}$ we execute Algorithm 1 that needs information only from systems $\Sigma_{[j]}$, $j \in \mathcal{N}_{M+1}$. If Algorithm 1 stops before the last step we declare that $\Sigma_{[M+1]}$ cannot be plugged in. Let $\mathcal{S}_i = \{j : i \in \mathcal{N}_j\}$ be the set of successors to system i . Since each system $\Sigma_{[j]}$, $j \in \mathcal{S}_{M+1}$ has the new neighbor $\Sigma_{[M+1]}$, it can be happen that existing matrices K_j , $j \in \mathcal{S}_{M+1}$ now give $\alpha_j \geq 1$ or $\hat{L}_j \leq 0$ or $\beta_i(\delta_i) \geq 1$. Indeed, when $\mathcal{N}_j$ gets larger, the quantity α_j in (19) (respectively $\hat{L}_j$ in (21)) can only increase (respectively decrease). Furthermore, the size of the set $\mathbb{Z}_j(\delta_j)$ increases and therefore the condition in (22) could be violated. This means

that for each $j \in \mathcal{S}_{M+1}$ the controllers $\mathcal{C}_{[j]}$ must be redesigned according to Algorithm 1. Again, if Algorithm 1 stops before completion for some $j \in \mathcal{S}_{M+1}$, we declare that $\Sigma_{[M+1]}$ cannot be plugged in.

In conclusion, the addition of system $\Sigma_{[M+1]}$ triggers the design of controller $\mathcal{C}_{[M+1]}$ and the redesign of controllers $\mathcal{C}_{[j]}$, $j \in \mathcal{S}_{M+1}$ according to Algorithm 1. Note that controller redesign does not propagate further in the network, i.e. even without changing controllers $\mathcal{C}_{[i]}$, $i \notin \{M+1\} \cup \mathcal{S}_{M+1}$ stability of the origin and constraint satisfaction are guaranteed for the new closed-loop system.

4.2 Unplugging operation

We consider the unplugging of system $\Sigma_{[k]}$, $k \in \mathcal{M}$. Since for each $i \in \mathcal{S}_k$ the set $\mathcal{N}_i$ gets smaller, we have that α_i in (19) (respectively $\hat{L}_i$ in (21)) cannot increase (respectively cannot decrease). Furthermore, the size of the set $\mathbb{Z}_i(\delta_i)$ cannot increase and therefore the inequality (22) cannot be violated. This means that for each $i \in \mathcal{S}_k$ the controller $\mathcal{C}_{[i]}$ does not have to be redesigned. Moreover since for each system $\Sigma_{[j]}$, $j \notin \{k\} \cup \mathcal{S}_k$ the set $\mathcal{N}_j$ does not change, the redesign of controller $\mathcal{C}_{[j]}$ is not required.

In conclusion, removal of system $\Sigma_{[k]}$ does not require the redesign of any controller, in order to guarantee stability of the origin and constraints satisfaction for the new closed-loop system. However we highlight that since systems $\Sigma_{[i]}$, $i \in \mathcal{S}_k$ have one neighbor less, the redesign of controllers $\mathcal{C}_{[i]}$ through Algorithm 1 could improve the performance.

5 Practical design and computational aspects

5.1 Automatic design of K_i and δ_i

The most difficult part of Algorithm 1 is step 1 and in this Section we propose an automatic method for computing the matrix K_i and $\delta_i > 0$. We assume that K_i is the LQ controller associated to matrices $Q_i \geq 0$ and $R_i > 0$, i.e.

$$K_i = (R_i + B_i^T \bar{P}_i B_i)^{-1} B_i^T \bar{P}_i A_{ii} \quad (24)$$

where $\bar{P}_i$ is the solution of the stationary Riccati equation

$$A_{ii}^T \bar{P}_i A_{ii} + Q_i - A_{ii}^T \bar{P}_i B_i (R_i + B_i^T \bar{P}_i B_i)^{-1} B_i^T \bar{P}_i A_{ii} = \bar{P}_i.$$

We then solve the following nonlinear optimization problem

$$\min_{\delta_i, Q_i, R_i} \mu_{\alpha_i} \alpha_i + \mu_{\beta_i} \beta_i(\delta_i) \quad (25a)$$

$$Q_i \geq 0, R_i > 0 \quad (25b)$$

$$\delta_i > 0 \quad (25c)$$

$$\alpha_i < 1, \hat{L}_i > 0 \quad (25d)$$

$$\beta_i(\delta_i) < 1 \quad (25e)$$

where $\mu_{\alpha_i} \geq 0$ and $\mu_{\beta_i} \geq 0$.

Feasibility of problem (25) guarantees that Algorithm 1 does not stop and then the controller $\mathcal{C}_{[i]}$ can be successfully designed. Moreover, in (25a) weights μ_{α_i} and μ_{β_i} establish a

trade-off between the maximization of sets $\hat{\mathbb{X}}_i$ and $\mathbb{V}_i$, respectively. A few remarks on the computations required for solving (25) are in order. First, beside the computation of K_i as in (24), problem (25) requires the computation of the set $\mathbb{Z}_i(\delta_i)$ that can be done using methods in [RKKM05], simplified as follows. Under Assumption 2, $\mathbb{W}_i = \oplus_{j \in \mathcal{N}_i} A_{ij} \mathbb{X}_j$ is a zonotope set defined as $\mathbb{W}_i = \{w_{[i]} = \Xi_{w_i} d_{w_i}, \|d_{w_i}\|_\infty \leq 1\}$. Hence, using the procedure proposed in [RKKM05], the set $\mathbb{Z}_i(\delta_i)$ is also a zonotope, defined as $\mathbb{Z}_i(\delta_i) = \{z_{[i]} = \Xi_{z_i} d_{z_i}, \|d_{z_i}\|_\infty \leq 1\}$, where $\Xi_{z_i} = \begin{bmatrix} \Xi_{w_i} & F_{ii} \Xi_{w_i} & \dots & F_{ii}^{s_i-1} \Xi_{w_i} \end{bmatrix}$ with s_i computed using Algorithm 1 in [RKKM05]. Since $\mathbb{W}_i$ and $\mathbb{Z}_i(\delta_i)$ are zonotopes, using (1), we can explicitly calculate the support function used in Algorithm 1 in [RKKM05] and rewrite (23) as

$$\hat{l}_{v_{i,r}}(\delta_i) = \|\Xi_{z_i}^T K_i^T h_{i,r}\|_1, \forall r \in 1 : r_{u_i}.$$

Second, we highlight that in absence of input constraints $\mathbb{U}_i$, constraint (25e) (and hence the computation of RPI sets $\mathbb{Z}_i(\delta_i)$) is not necessary. Indeed if $\mathbb{U}_i = \mathbb{R}^{m_i}$, the inclusion (16) holds for all sets $\mathbb{V}_i \subseteq \mathbb{R}^{m_i}$. Third, the series in (19) and (20) involve only positive terms and can be easily truncated either if (25d) is violated or if summands fall below the machine precision. Finally, in order to simplify problem (25) one can assume $Q_i = \text{diag}(q_{i,1}, \dots, q_{i,n_i})$ and $R_i = \text{diag}(r_{i,1}, \dots, r_{i,m_i})$ hence replacing the matrix inequalities in (25b) with the scalar inequalities $q_{i,k} \geq 0$, $k \in 1 : n_i$ and $r_{i,k} > 0$, $k \in 1 : m_i$.

5.2 Parameter-dependent subsystem

In many engineering applications parameters of subsystem i are influenced by neighboring subsystems. We model this scenario replacing (4) with

$$\Sigma_{[i]}^p : \quad x^+_{[i]} = A_{ii}(\xi_{ii}, \{\xi_{ij}\}_{j \in \mathcal{N}_i}) x_{[i]} + B_i(\xi_{ii}, \{\xi_{ij}\}_{j \in \mathcal{N}_i}) u_{[i]} + \sum_{j \in \mathcal{N}_i} A_{ij} x_{[j]} \quad (26)$$

where $\xi_{ij} \in \mathbb{R}^{p_{ij}}$ are parameter vectors.

We highlight that for a given sets $\mathcal{N}_i$, $i \in \mathcal{M}$, matrices A_{ii} and B_i are constant and design of P&P DeMPC regulators can be still done using the methods described in Section 3. Furthermore, the procedure for plugging in a new system discussed in Section 4.1 can be applied with no change since it requires the redesign of controllers $\mathcal{C}_{[j]}$, $j \in \mathcal{S}_{M+1}$, i.e. controllers associated to the subsystems $\Sigma_{[j]}^p$ for which matrices A_{jj} and B_j could change. However, when system $\Sigma_{[k]}^p$ gets unplugged, it is now mandatory to retune all controllers $\mathcal{C}_{[j]}$, $j \in \mathcal{S}_k$ since changes of matrices A_{jj} and B_j could hamper the fulfillment of conditions (19), (21) or (22) when using the matrices K_j and the scalars δ_j computed prior to the subsystem removal. Moreover, if Algorithm 1 stops before completing the redesign of controllers $\mathcal{C}_{[j]}$, $\forall j \in \mathcal{S}_k$, we declare that subsystem $\Sigma_{[k]}^p$ cannot be unplugged.

6 Example: Power Network System

In this Section, we apply the proposed DeMPC scheme to a power network system composed by several power generation areas coupled through tie-lines. We aim at designing the AGC layer with the goals of

- keeping the frequency approximately at a nominal value;

- controlling the tie-line powers in order to reduce power exchanges between areas. In the asymptotic regime each area should compensate for local load steps and produce the required power.

In particular we will show advantages brought about by P&P DeMPC when generation areas are connected/disconnected to/from an existing network.

The dynamics of an area equipped with primary control and linearized around equilibrium value for all variables can be described by the following continuous-time LTI model [Saa02]

$$\Sigma_{[i]}^C : \quad \dot{x}_{[i]} = A_{ii}x_{[i]} + B_i u_{[i]} + L_i \Delta P_{L_i} + \sum_{j \in \mathcal{N}_i} A_{ij} x_{[j]} \quad (27)$$

where $x_{[i]} = (\Delta\theta_i, \Delta\omega_i, \Delta P_{m_i}, \Delta P_{v_i})$ is the state, $u_{[i]} = \Delta P_{ref_i}$ is the control input of each area, ΔP_L is the local power load and $\mathcal{N}_i$ is the sets of neighboring areas, i.e. areas directly connected to $\Sigma_{[i]}^C$ through tie-lines. The matrices of system (27) are defined as

$$\begin{aligned} A_{ii}(\{P_{ij}\}_{j \in \mathcal{N}_i}) &= \begin{bmatrix} 0 & 1 & 0 & 0 \\ -\frac{\sum_{j \in \mathcal{N}_i} P_{ij}}{2H_i} & -\frac{D_i}{2H_i} & \frac{1}{2H_i} & 0 \\ 0 & 0 & -\frac{1}{T_{t_i}} & \frac{1}{T_{t_i}} \\ 0 & -\frac{1}{R_i T_{g_i}} & 0 & -\frac{1}{T_{g_i}} \end{bmatrix} & B_i = \begin{bmatrix} 0 \\ 0 \\ 0 \\ \frac{1}{T_{g_i}} \end{bmatrix} \\ A_{ij} &= \begin{bmatrix} 0 & 0 & 0 & 0 \\ \frac{P_{ij}}{2H_i} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} & L_i = \begin{bmatrix} 0 \\ -\frac{1}{2H_i} \\ 0 \\ 0 \end{bmatrix} \end{aligned} \quad (28)$$

For the meaning of constants as well as parameter values we defer the reader to Appendix C. We highlight that all parameter values are within the range of those used in Chapter 12 of [Saa02].

We note that model (27) is input decoupled since both ΔP_{ref_i} and ΔP_{L_i} act only on subsystem $\Sigma_{[i]}^C$. Moreover, subsystems $\Sigma_{[i]}^C$ are parameter dependent since the local dynamics depends on the quantities $-\frac{\sum_{j \in \mathcal{N}_i} P_{ij}}{2H_i}$. We equip each subsystem $\Sigma_{[i]}^C$ with the constraints on $\Delta\theta_i$ and on ΔP_{ref_i} specified in Appendix C. We obtain models $\Sigma_{[i]}$ by discretizing models $\Sigma_{[i]}^C$ with 1 sec sampling time, using exact discretization and treating $u_{[i]}$, ΔP_{L_i} , $x_{[j]}$, $j \in \mathcal{N}_i$ as exogenous signals.

In the following we first design the AGC layer for a power network composed by four areas (Scenario 1) and then we show how in presence of connection/disconnection of an area (Scenario 2 and 3, respectively) the AGC can be redesigned via plugging in and unplugging operations.

6.1 Scenario 1

We consider four areas interconnected as in Figure 1. For each system $\Sigma_{[i]}$ we synthesize the controller K_i , $i \in \mathcal{M}$ solving an LQ problem for the nominal system, as shown in Section 5.1 with $\mu_{\alpha_i} = 1$ and $\mu_{\beta_i} = 1$, $\forall i \in \mathcal{M}$, and obtain the following matrices

$$\begin{aligned} K_1 &= -\begin{bmatrix} 0.508 & 0.201 & 0.006 & 0.001 \end{bmatrix}, & K_2 &= -\begin{bmatrix} 0.729 & 0.437 & 0.008 & 0.002 \end{bmatrix}, \\ K_3 &= -\begin{bmatrix} 3.409 & 4.759 & 0.090 & 0.030 \end{bmatrix}, & K_4 &= -\begin{bmatrix} 4.426 & 6.348 & 0.233 & 0.038 \end{bmatrix}, \end{aligned} \quad (29)$$

that allow inequalities (19) to be fulfilled. Hence $\mathbf{K}$ verifies Assumption 1-(ii). Setting $\delta_i = 10^{-4}$, $\forall i \in \mathcal{M}$ and applying steps 2-5 of Algorithm 1, we can compute sets $\mathbb{Z}_i(\delta_i)$, $\hat{\mathbb{X}}_i$ and $\mathbb{V}_i$

Figure 1: Power network system of Scenario 1

such that inclusions (15) and (16) hold, $\forall i \in \mathcal{M}$. Control variables $u_{[i]}$ are obtained through (6) where $v_{[i]} = \kappa_i(x_{[i]})$ and $\bar{x}_{[i]} = \eta_i(x_{[i]})$ are computed at each time t solving the optimization problem (18) and replacing the cost function in (18a) with the following one depending upon $x_{[i]}^O = (0, 0, \Delta P_{L_i}, \Delta P_{L_i})$ and $u_{[i]}^O = \Delta P_{L_i}$

$$J_i^{N_i} = \sum_{k=t}^{t+N_i-1} (\|\hat{x}_{[i]}(k) - x_{[i]}^O\|_{\hat{Q}_i} + \|v_{[i]}(k) - u_{[i]}^O\|_{\hat{R}_i} + \|x_{[i]}(t+N_i) - x_{[i]}^O\|_{\hat{S}_i}). \quad (30)$$

Note that, except for the above modification of the cost function, that is needed for counteracting load disturbances, we followed exactly the design procedure described in Section 2. Moreover, we highlight that each area can locally absorb the load steps specified in Table 3 of Appendix C. This is also shown by convergence to zero of the power transfer between areas i and j given by

$$\Delta P_{tie_{ij}} = P_{ij}(\Delta\theta_i - \Delta\theta_j) \quad (31)$$

and represented in Figure 3.

In Figure 2 we compare the performance of proposed DeMPC with the performance of centralized MPC. For centralized MPC we consider the overall system composed by the four areas, use the cost function $\sum_{i \in \mathcal{M}} J_i^N$ and impose the collective constraints (8). The prediction horizon is $N_i = 20$, $i \in \mathcal{M}$ for MPC- i controllers and $N = 20$ for centralized MPC. In the control experiment, step power loads ΔP_{L_i} specified in Appendix C have been used and they account for the step-like changes of the control variables in Figure 2. We highlight that the performance of decentralized and centralized MPC are totally comparable, in terms of frequency deviation (Figure 2(a)), control variables (Figure 2(b)) and power transfers $\Delta P_{tie_{ij}}$ (Figure 3).

6.2 Scenario 2

We consider the power network proposed in Scenario 1 and we add a fifth area connected as in Figure 4 with values of parameters and constraints listed in Table 2 of Appendix C. Therefore, the set of successors to system 5 is $\mathcal{S}_5 = \{2, 4\}$.

As described in Section 4.1, only systems $\Sigma_{[j]}$, $j \in \mathcal{S}_5$ update their controller $\mathcal{C}_{[j]}$. For systems $\Sigma_{[j]}$, $j \in \mathcal{S}_5$, since the set $\mathcal{N}_j$ changes, we retune controllers $\mathcal{C}_{[j]}$ using Algorithm 1. In particular, we compute K_j , $j \in \mathcal{S}_5$ and K_5 using the procedure described in Section 5.1 with $\mu_{\alpha_k} = 1$ and

(a) Frequency deviation in each area controlled by the proposed De-MPC (bold line) and centralized MPC (dashed line).

(b) Load reference set-point in each area controlled by the proposed De-MPC (bold line) and centralized MPC (dashed line).

Figure 2: Simulation Scenario 1: **2(a)** Frequency deviation and **2(b)** Load reference in each area.

$\mu_{\beta_k} = 1, k \in \{5\} \cup \mathcal{S}_5$ and obtain

$$\begin{aligned} K_2 &= - \begin{bmatrix} 0.659 & 1.275 & 0.028 & 0.007 \end{bmatrix}, \quad K_4 = - \begin{bmatrix} 0.713 & 1.105 & 0.048 & 0.008 \end{bmatrix}, \\ K_5 &= - \begin{bmatrix} 0.123 & 0.158 & 0.007 & 0.001 \end{bmatrix}, \end{aligned} \quad (32)$$

that allow inequalities (19) to be verified for systems $\Sigma_{[j]}, j \in \mathcal{S}_5$ and $\Sigma_{[5]}$. Therefore $\mathbf{K}$ fulfills Assumption 1-(ii). Setting $\delta_j = 10^{-4}, j \in \mathcal{S}_5$ and $\delta_5 = 10^{-4}$, the execution of Algorithm 1 does not stop before completion and hence we compute the new sets $\mathbb{Z}_j(\delta_j), \hat{\mathbb{X}}_j$ and $\mathbb{V}_j, j \in \{5\} \cup \mathcal{S}_5$.

Figure 3: Simulation Scenario 1: tie-line power between each area controlled by the proposed DeMPC (bold line) and centralized MPC (dashed line).

Figure 4: Power network system of Scenario 2

We highlight that no retuning of controllers $\mathcal{C}_{[1]}$ and $\mathcal{C}_{[3]}$ is needed since systems $\Sigma_{[1]}$ and $\Sigma_{[3]}$ are not neighbors to system $\Sigma_{[5]}$.

In Figure 5 we compare the performance of proposed DeMPC with the performance of centralized MPC. For centralized MPC we consider the overall system composed by the four areas, use the cost function $\sum_{i \in \mathcal{M}} J_i^N$ and impose the collective constraints (8). The prediction horizon is $N_i = 20$, $i \in \mathcal{M}$ for MPC- i controllers and $N = 20$ for centralized MPC. In the control experiment, step power loads ΔP_{L_i} specified in Appendix C have been used and they account for the step-like changes of the control variables in Figure 5. We highlight that the performance of decentralized and centralized MPC are totally comparable, in terms of frequency deviation (Figure 5(a)), control variables (Figure 5(b)) and power transfers $\Delta P_{tie_{ij}}$ (Figure 6).

(a) Frequency deviation in each area controlled by the proposed De-MPC (bold line) and centralized MPC (dashed line).

(b) Load reference set-point in each area controlled by the proposed De-MPC (bold line) and centralized MPC (dashed line).

Figure 5: Simulation Scenario 2: **5(a)** Frequency deviation and **5(b)** Load reference in each area.

6.3 Scenario 3

We consider the power network described in Scenario 2 and disconnect the area 4, hence obtaining the areas connected as in Figure 7. The set of successors to system 4 is $\mathcal{S}_4 = \{3, 5\}$. Because of disconnection, systems $\Sigma_{[j]}^C$, $j \in \mathcal{S}_4$ change their neighbors and local dynamics A_{jj} . Moreover, it is possible to verify that matrices K_j computed in Scenario 2 do not solve Problem $\mathcal{P}_j$, $j \in \mathcal{S}_4$. Then as described in Section 5.2, each subsystem $\Sigma_{[j]}^C$, $j \in \mathcal{S}_4$ must retune controller $\mathcal{C}_{[j]}$ by running Algorithm 1. In particular, we compute K_3 and K_5 using the procedure proposed in Section 5.1

Figure 6: Simulation Scenario 2: tie-line power between each area controlled by the proposed DeMPC (bold line) and centralized MPC (dashed line).

Figure 7: Power network system of Scenario 3

with $\mu_{\alpha_j} = 1$ and $\mu_{\beta_j} = 1$, $j \in \mathcal{S}_4$ and obtain

$$K_3 = - \begin{bmatrix} 4.766 & 5.954 & 0.110 & 0.036 \end{bmatrix}, \quad K_5 = - \begin{bmatrix} 4.102 & 4.861 & 0.201 & 0.038 \end{bmatrix}, \quad (33)$$

that allows one to verify inequalities (19) for systems $\Sigma_{[j]}$, $j \in \mathcal{S}_4$. Therefore $\mathbf{K}$ is such that Assumption 1-(ii) holds. Setting $\delta_j = 10^{-4}$, $j \in \mathcal{S}_4$, the execution of Algorithm 1 does not stop before completion and hence we compute the new sets $\mathbb{Z}_j(\delta_j)$, $\hat{\mathbb{X}}_j$ and $\mathbb{V}_j$, $j \in \mathcal{S}_4$. We highlight that retuning of controllers $\mathcal{C}_{[1]}$ and $\mathcal{C}_{[2]}$ is not needed since systems $\Sigma_{[1]}$ and $\Sigma_{[2]}$ are not neighbors to system $\Sigma_{[4]}$.

In Figure 8 we compare the performance of proposed DeMPC with the performance of centralized MPC. For centralized MPC we consider the overall system composed by the four areas, use the cost function $\sum_{i \in \mathcal{M}} J_i^N$ and impose the collective constraints (8). The prediction horizon is $N_i = 20$, $i \in \mathcal{M}$ for MPC- i controllers and $N = 20$ for centralized MPC. In the control ex-

(a) Frequency deviation in each area controlled by the proposed De-MPC (bold line) and centralized MPC (dashed line).

(b) Load reference set-point in each area controlled by the proposed De-MPC (bold line) and centralized MPC (dashed line).

Figure 8: Simulation Scenario 3: **8(a)** Frequency deviation and **8(b)** Load reference in each area.

periment, step power loads ΔP_{L_i} specified in Appendix C have been used also in this case. We highlight that the performance of decentralized and centralized MPC are totally comparable in terms of frequency deviation (Figure 8(a)), control variables (Figure 8(b)) and power transfers $\Delta P_{tie_{ij}}$ (Figure 9).

Figure 9: Simulation Scenario 3: tie-line power between each area controlled by the proposed DeMPC (bold line) and centralized MPC (dashed line).

7 Conclusions

In this paper we proposed a tube-based DeMPC scheme for linear constrained systems, with the goal of stabilizing the origin of the closed-loop system and guaranteeing constraints satisfaction. The key feature of our approach is that the design procedure does not require any centralized computation. This enables P&P operations, i.e. when a subsystem is plugged-in or unplugged at most the synthesis of its controller and the redesign of successors' controllers are needed. In future we will generalize our approach to embrace decentralized output-feedback control and tracking problems.

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A Proof of Theorem 1

The proof uses arguments similar to the ones adopted in [FS12] for proving Theorem 1.

We first show recursive feasibility, i.e. that $x_{[i]}(t) \in \mathbb{X}_i^N$, $\forall i \in \mathcal{M}$ implies $x_{[i]}(t+1) \in \mathbb{X}_i^N$.

Assume that, at instant t , $x_{[i]}(t) \in \mathbb{X}_i^N$. The optimal nominal input and state sequences obtained by solving each MPC- i problem $\mathbb{P}_i^N$ are $v_{[i]}(0 : N_i - 1|t) = \{v_{[i]}(0|t), \dots, v_{[i]}(N_i - 1|t)\}$ and $\hat{x}_{[i]}(0 : N_i|t) = \{\hat{x}_{[i]}(0|t), \dots, \hat{x}_{[i]}(N_i|t)\}$, respectively. Define $v_{[i]}^{aux}(N_i|t) = \kappa_i^{aux}(\hat{x}_{[i]}(N_i|t))$ and compute $\hat{x}_{[i]}^{aux}(N_i + 1|t)$ according to (18c) from $\hat{x}_{[i]}(N_i|t)$ and $v_{[i]}(N_i|t) = v_{[i]}^{aux}(N_i|t)$. Note that, in view of constraint (18f) and points (ii) and (iii) of Assumption 4, $v_{[i]}^{aux}(N_i|t) \in \mathbb{V}_i$ and $\hat{x}_{[i]}^{aux}(N_i + 1|t) \in \hat{\mathbb{X}}_{f_i} \subseteq \tilde{\mathbb{X}}_i$. We also define the input sequence

$$\bar{v}_{[i]}(1 : N_i|t) = \{v_{[i]}(1|t), \dots, v_{[i]}(N_i - 1|t), v_{[i]}^{aux}(N_i|t)\} \quad (34)$$

and the state sequence produced by (18c) from the initial condition $\hat{x}_{[i]}(0|t)$ and the input sequence $\bar{v}_{[i]}(1 : N_i|t)$, i.e.

$$\tilde{\hat{x}}_{[i]}(1 : N_i + 1|t + 1) = \{\hat{x}_{[i]}(1|t), \dots, \hat{x}_{[i]}(N_i|t), \hat{x}_{[i]}^{aux}(N_i + 1|t)\}. \quad (35)$$

In view of the constraints (18) at time t and recalling that $\mathbb{Z}_i(\delta_i)$ is an RPI for (18) and $w_{[i]} \in \mathbb{W}_i = \bigoplus_{j \in \mathcal{N}_i} A_{ij} \mathbb{X}_j$, we have that $x_{[i]}(t+1) - \hat{x}_{[i]}(1|t) \in \mathbb{Z}_i(\delta_i)$. Therefore, we can conclude that the state and the input sequences $\tilde{\hat{x}}_{[i]}(1 : N_i + 1|t)$ and $\bar{v}_{[i]}(1 : N_i|t)$ are feasible at $t+1$, since constraints (18b)-(18f) are satisfied. This proves recursive feasibility.

We now prove convergence of the optimal cost function.

We define $\mathbb{P}_i^{N,0}(\hat{x}_{[i]}(0|t)) = \min_{v_{[i]}(0:N_i-1|t)} \sum_{k=0}^{N_i-1} \ell_i(\hat{x}_{[i]}(k), v_{[i]}(k)) + V_{f_i}(\hat{x}_{[i]}(N_i))$ subject to the constraints (18c)-(18f). By optimality, using the feasible control law (34) and the corresponding state sequence (35) one has

$$\mathbb{P}_i^{N,0}(\hat{x}_{[i]}(1|t)) \leq \sum_{k=1}^{N_i} \ell_i(\hat{x}_{[i]}(k|t), v_{[i]}(k|t)) + V_{f_i}(\hat{x}_{[i]}^{aux}(N_i + 1|t + 1)) \quad (36)$$

where it has been set $v_{[i]}(N_i|t) = v_{[i]}^{aux}(N_i|t)$. Therefore we have

$$\begin{aligned} \mathbb{P}_i^{N,0}(\hat{x}_{[i]}(1|t)) - \mathbb{P}_i^{N,0}(\hat{x}_{[i]}(0|t)) &\leq -\ell_i(\hat{x}_{[i]}(0|t), v_{[i]}(0|t)) + \ell_i(\hat{x}_{[i]}(N_i|t), v_{[i]}^{aux}(N_i|t)) + \\ &\quad + V_{f_i}(\hat{x}_{[i]}^{aux}(N_i + 1|t)) - V_{f_i}(\hat{x}_{[i]}^{aux}(N_i|t)). \end{aligned} \quad (37)$$

In view of Assumption 4-(iv), from (37) we obtain

$$\mathbb{P}_i^{N,0}(\hat{x}_{[i]}(1|t)) - \mathbb{P}_i^{N,0}(\hat{x}_{[i]}(t)) \leq -\ell_i(\hat{x}_{[i]}(t), v_{[i]}(t)) \quad (38)$$

and therefore $\hat{x}_{[i]}(0|t) \rightarrow 0$ and $v_{[i]}(0|t) \rightarrow 0$ as $t \rightarrow \infty$.

Next we prove convergence to zero of state trajectories $\mathbf{x}(t)$ of the closed-loop system with $\mathbf{x}(0) \in \mathbb{X}^N$.

Recall that the state $\mathbf{x}(t)$ evolves according to the equation (7). By asymptotic convergence to zero of the nominal state and input signals $\hat{x}_{[i]}(0|t)$ and $v_{[i]}(0|t)$ respectively, using the diagonal structure of $\mathbf{B}$ and $\mathbf{K}$, we obtain that $\mathbf{B}(\mathbf{v}(0|t) - \mathbf{K}\hat{\mathbf{x}}(0|t))$ is an asymptotically vanishing term. Under Assumption 1-(ii), $\mathbf{A} + \mathbf{BK}$ is Schur, hence we obtain $\mathbf{x}(t) \rightarrow 0$ as $t \rightarrow \infty$.

We prove now stability of the origin of the closed-loop system therefore completing the proof of statement (ii). We first show that

$$\mathbf{x}(0) \in \mathbb{Z} \Rightarrow \mathbf{x}(t) \in \mathbb{Z} \text{ and } \mathbf{x}(t+1) = (\mathbf{A} + \mathbf{BK})\mathbf{x}(t), \forall t \geq 0 \quad (39)$$

where $\mathbb{Z} = \prod_{i \in \mathcal{M}} \mathbb{Z}_i(\delta_i)$. Formula (39) is an easy consequence of Proposition 2 in [MSR05] that we detail in the following for the sake of completeness.

If $\mathbf{x}(0) \in \mathbb{Z}$ then, as shown in [MSR05], $\mathbf{v} = 0$ and $\bar{\mathbf{x}} = 0$. Therefore from (7) we have

$$\mathbf{x}^+(1) = (\mathbf{A} + \mathbf{BK})\mathbf{x}(0)$$

Furthermore, since sets $\mathbb{Z}_i$ are RPI for (14), one has that $\mathbb{Z}$ is positively invariant for

$$\mathbf{x}^+ = (\mathbf{A} + \mathbf{BK})\mathbf{x} \quad (40)$$

that coincides with (14) after renaming variables $z_{[i]}$ as $x_{[i]}$. Therefore $\mathbf{x}(t) \in \mathbb{Z}$ for $t = 1$ and, applying the previous argument recursively for $t \geq 1$, we have shown that (39) holds.

Now, we focus on stability. Given $\epsilon > 0$, choose $\eta \in (0, 1)$ such that

$$\eta\mathbb{Z} \subset B_\epsilon(0). \quad (41)$$

Such an η always exists because $\mathbb{Z}$ is bounded and includes the origin in its interior. More precisely, boundedness of $\mathbb{Z}$ follows from $\mathbb{Z} \subset \mathbb{X}$ and the boundedness of $\mathbb{X}$, that is guaranteed by Assumption 2. Furthermore, the mRPI for (14) is given by [RM09]

$$\underline{\mathbb{Z}}_i = \bigoplus_{k=0}^{\infty} F_i^k \mathbb{W}_i$$

and therefore it includes the origin in its interior. It follows that the same is true for sets $\mathbb{Z}_i(\delta_i)$, $i \in \mathcal{M}$ and $\mathbb{Z} = \prod_{i \in \mathcal{M}} \mathbb{Z}_i(\delta_i)$. Since the origin is strictly contained in $\mathbb{Z}$, there always exists $\underline{\delta} > 0$ such that $B_{\underline{\delta}}(0) \subseteq \eta\mathbb{Z}$. Since $B_{\underline{\delta}}(0) \subset \mathbb{Z}$ one has that, in view of (39), the state trajectory $\mathbf{x}(t)$, $t \geq 0$ stemming from $\mathbf{x}(0) \in B_{\underline{\delta}}(0)$ fulfill the dynamics (40). Furthermore, since (40) is a linear system for which $\mathbb{Z}$ is positively invariant set, one has that also $\eta\mathbb{Z}$ is positively invariant. Then, we have shown that

$$\mathbf{x}(0) \in B_{\underline{\delta}}(0) \rightarrow \mathbf{x}(t) \in \eta\mathbb{Z}, \quad \forall t \geq 0.$$

From (41) stability of the origin follows. ■

B Proof of Theorem 2

B.1 Proof of (I)

Define a matrix $\mathbf{M}$ such that its ij -th entry μ_{ij} is

$$\begin{aligned} \mu_{ij} &= -1 && \text{if } i = j \\ \mu_{ij} &= \sum_{k=0}^{\infty} \|\mathcal{F}_i F_i^k A_{ij} \mathcal{F}_j^b\|_{\infty} && \text{if } i \neq j. \end{aligned}$$

Note that all the off-diagonal entries of matrix $\mathbf{M}$ are non-negative, i.e., it is Metzler [FR00]. We recall the following results.

Lemma 1 (see [MS07]). *Let matrix $\mathbf{M} \in \mathbb{R}^{M \times M}$ be Metzler. Then $\mathbf{M}$ is Hurwitz if and only if there is a vector $\nu \in \mathbb{R}_+^M$ such that $\mathbf{M}\nu < 0$.*

Lemma 2. *Define the matrix $\Gamma = \mathbf{M} + \mathbb{I}_M$ where $\mathbf{M} \in \mathbb{R}^{M \times M}$, $\mathbb{I}_M$ is the $M \times M$ identity matrix and Γ is non negative. Then the Metzler matrix $\mathbf{M}$ is Hurwitz if and only if Γ is Schur.*

The proof of Lemma 2 easily follows from Theorem 13 in [FR00]. Inequalities (19) are equivalent to $\mathbf{M}\nu < \mathbf{0}_M$ where $\nu = \mathbf{1}_M$. Then, from Lemma 1, $\mathbf{M}$ is Hurwitz. From Lemma 2, (19) implies that matrix $\Gamma = \mathbf{M} + \mathbb{I}_M$ is Schur. For system $\Sigma_{[i]}$ in (4)-(5), when $u_{[i]}$ is defined as in (6), $v_{[i]} = 0$ and $\bar{x}_{[i]} = 0$, we have

$$x_{[i]}(t) = F_i^t x_{[i]}(0) + \sum_{k=0}^{t-1} F_i^k \sum_{j \in \mathcal{N}_i} A_{ij} x_{[j]}(t-k-1) \quad (42)$$

In view of (42) we can write

$$\begin{aligned} \|\mathcal{F}_i x_{[i]}(t)\|_\infty &\leq \|\mathcal{F}_i F_i^t \mathcal{F}_i^\flat\|_\infty \|\mathcal{F}_i x_{[i]}(0)\|_\infty + \\ &\quad + \sum_{j \in \mathcal{N}_i} \gamma_{ij} \max_{k \leq t} \|\mathcal{F}_j x_{[j]}(k)\|_\infty. \end{aligned}$$

where γ_{ij} are the entries of Γ . Denoting $\tilde{x}_{[i]} = \mathcal{F}_i x_{[i]}$, we can collectively define $\tilde{\mathbf{x}} = \tilde{\mathcal{F}}\mathbf{x}$, where $\tilde{\mathcal{F}} = \text{diag}(\mathcal{F}_1, \dots, \mathcal{F}_M)$. From the definition of sets $\mathbb{X}_i$, we have $\text{rank}(\tilde{\mathcal{F}}) = n$. We define the system

$$\tilde{\mathbf{x}}^+ = (\tilde{\mathbf{A}} + \tilde{\mathbf{B}}\tilde{\mathbf{K}})\tilde{\mathbf{x}} \quad (43)$$

where $\tilde{\mathbf{A}} = \tilde{\mathcal{F}}\mathbf{A}\tilde{\mathcal{F}}^\flat$, $\tilde{\mathbf{B}} = \tilde{\mathcal{F}}\mathbf{B}$ and $\tilde{\mathbf{K}} = \mathbf{K}\tilde{\mathcal{F}}^\flat$. In order to analyze the stability of the origin of (43), we consider the method proposed in [DRW07]. In view of Corollary 16 in [DRW07], the overall system (43) is asymptotically stable if the gain matrix Γ is Schur. As shown above this property is implied by (19).

Moreover, system (43) is an expansion of the original system (see Chapter 3.4 in [Lun92]). In view of the inclusion principle (see Theorem 3.3 in [Lun92] and [Sta04] for a discrete-time version), the asymptotic stability of (43) implies the asymptotic stability of the original system.

B.2 Proof of (II)

First note that, for $i \in \mathcal{M}$, in view of (10) $\|f_{i,r}^T \Xi_i\|_\infty = 1$ for all $r \in 1 : \bar{r}_i$ and therefore $\|\mathcal{F}_i \Xi_i\|_\infty = 1$. This implies that $\|f_{i,r}^T F_i^k A_{ij} \Xi_j\|_\infty \leq \|f_{i,r}^T F_i^k A_{ij} \mathcal{F}_j^\flat\|_\infty \|\mathcal{F}_j \Xi_j\|_\infty = \|f_{i,r}^T F_i^k A_{ij} \mathcal{F}_j^\flat\|_\infty \leq \|\mathcal{F}_i F_i^k A_{ij} \mathcal{F}_j^\flat\|_\infty$. Therefore, in view of (19), for all $r \in 1 : \bar{r}_i$

$$\sum_{k=0}^{\infty} \sum_{j \in \mathcal{N}_i} \|f_{i,r} F_i^k A_{ij} \Xi_j\|_\infty \leq \sum_{k=0}^{\infty} \sum_{j \in \mathcal{N}_i} \|\mathcal{F}_i F_i^k A_{ij} \mathcal{F}_j^\flat\|_\infty < 1 \quad (44)$$

Now we want to find parameter $\hat{l}_i > 0$ such that, simultaneously, the inclusion (15) holds and $\mathbb{Z}_i(\delta_i)$ is a δ_i -outer approximation of the mRPI $\underline{\mathbb{Z}}_i$. The mRPI for (14) is given by [RKKM05]

$$\underline{\mathbb{Z}}_i = \bigoplus_{k=0}^{\infty} F_i^k \bigoplus_{j \in \mathcal{N}_i} A_{ij} \mathbb{X}_j. \quad (45)$$

From [RKKM05], for given $\delta_i > 0$ there exist $\alpha_i \in \mathbb{R}$ and $s_i \in \mathbb{N}_+$ such that the set

$$\mathbb{Z}_i(\delta_i) = (1 - \alpha_i)^{-1} \bigoplus_{k=0}^{s_i-1} F_i^k \bigoplus_{j \in \mathcal{N}_i} A_{ij} \mathbb{X}_j \quad (46)$$

is a δ_i -outer approximation of the mRPI $\mathbb{Z}_i$.

Define $\bar{\mathbb{X}}_i = \hat{\mathbb{X}}_i \oplus \mathbb{Z}_i(\delta_i)$. Following the proof of Proposition 2 in [FS12] and using arguments from Section 3 of [KG98], we can then guarantee (15) if $\bar{\mathbb{X}}_i \subseteq \mathbb{X}_i$, which holds if, for all $r \in 1 : \bar{r}_i$

$$\sup_{\substack{z_{[i]} \in \mathbb{Z}_i(\delta_i) \\ \hat{x}_{[i]} \in \hat{\mathbb{X}}_i}} f_{i,r}^T(z_{[i]} + \hat{x}_{[i]}) \leq 1. \quad (47)$$

Using (2) and (45), the inequalities (47) are verified if

$$\sup_{\substack{\{x_{[j]}(k) \in \mathbb{X}_j\}_{j \in \mathcal{N}_i}^{k=0, \dots, \infty} \\ \hat{x}_{[i]} \in \hat{\mathbb{X}}_i \\ \sigma_i \in B_{\delta_i}(0)}} h_{i,r}^x(\{x_{[j]}(k)\}_{j \in \mathcal{N}_i}^{k=0, \dots, \infty}, \hat{x}_{[i]}) + \|f_{i,r}^T \sigma_i\|_\infty \leq 1 \quad (48)$$

where $h_{i,r}^x(\cdot) = f_{i,r}^T(\sum_{k=0}^{\infty} F_i^k \sum_{j \in \mathcal{N}_i} A_{ij} x_{[j]}(k) + \hat{x}_{[i]})$.

Since $\|f_{i,r}^T \sigma_i\|_\infty \leq \|f_{i,r}^T\|_\infty \delta_i$, conditions (48) are satisfied if

$$\sup_{\substack{\{x_{[j]}(k) \in \mathbb{X}_j\}_{j \in \mathcal{N}_i}^{k=0, \dots, \infty} \\ \hat{x}_{[i]} \in \hat{\mathbb{X}}_i}} h_{i,r}^x(\{x_{[j]}(k)\}_{j \in \mathcal{N}_i}^{k=0, \dots, \infty}, \hat{x}_{[i]}) \leq 1 - \|f_{i,r}^T\|_\infty \delta_i. \quad (49)$$

Using (10) and (11) we can rewrite (49) as

$$\sup_{\substack{\{\|d_j(k)\|_\infty \leq 1\}_{j \in \mathcal{N}_i}^{k=0, \dots, \infty} \\ \|\hat{d}_i\|_\infty \leq \hat{l}_i}} h_{i,r}^d(\{d_j(k)\}_{j \in \mathcal{N}_i}^{k=0, \dots, \infty}, \hat{d}_i) \leq 1 - \|f_{i,r}^T\|_\infty \delta_i \quad (50)$$

where $h_{i,r}^d(\cdot) = f_{i,r}^T(\sum_{k=0}^{\infty} F_i^k \sum_{j \in \mathcal{N}_i} A_{ij} \Xi_j d_j(k) + \hat{\Xi}_i \hat{d}_i)$.

The inequalities (50) are satisfied if

$$\sum_{k=0}^{\infty} \sum_{j \in \mathcal{N}_i} \|f_{i,r}^T F_i^k A_{ij} \Xi_j\|_\infty + \|f_{i,r}^T \hat{\Xi}_i\|_\infty \hat{l}_i \leq 1 - \|f_{i,r}^T\|_\infty \delta_i \quad (51)$$

for all $r \in 1 : \bar{r}_i$.

In view of (44) there exist sufficiently small $\delta_i > 0$ and $\hat{l}_i > 0$ satisfying (51) (and therefore verifying (15)), e.g. choosing $\hat{l}_i \in (0, \hat{L}_i]$.

B.3 Proof of (III)

For each $i \in \mathcal{M}$, we want to find tightened input constraint $\mathbb{V}_i$ such that (16) holds. Following the rational used in Section 3 of [KG98], from definition of sets $\mathbb{U}_i$ and $\mathbb{V}_i$, (16) holds if (22) is satisfied. Hence, choosing $\mathbb{V}_i$ as in (13), for $l_{v_{i,r}} = \hat{l}_{v_{i,r}}(\delta_i)$ the inclusion (16) holds. ■

C Parameters, constraints and setpoints of experiment described in Section 6

$\Delta\theta_i$	Deviation of the angular displacement of the rotor with respect to the stationary reference axis on the stator
$\Delta\omega_i$	Speed deviation of rotating mass from nominal value
ΔP_{m_i}	Deviation of the mechanical power from nominal value (p.u.)
ΔP_{v_i}	Deviation of the steam valve position from nominal value (p.u.)
ΔP_{ref_i}	Deviation of the reference set power from nominal value (p.u.)
ΔP_{L_i}	Deviation of the nonfrequency-sensitive load change from nominal value (p.u.)
H_i	Inertia constant defined as $H_i = \frac{\text{kinetic energy at rated speed}}{\text{machine rating}}$ (typically values in range [1 – 10] sec)
R_i	Speed regulation
D_i	Defined as $\frac{\text{percent change in load}}{\text{change in frequency}}$
T_{t_i}	Prime mover time constant (typically values in range [0.2 – 2] sec)
T_{g_i}	Governor time constant (typically values in range [0.1 – 0.6] sec)
P_{ij}	Slope of the power angle curve at the initial operating angle between area i and area j

Table 1: Variables of a generation area with typical value ranges [Saa02]. (p.u.) stands for “per unit”.

	Area 1	Area 2	Area 3	Area 4	Area 5
H_i	12	10	8	8	10
R_i	0.05	0.0625	0.08	0.08	0.05
D_i	0.7	0.9	0.9	0.7	0.86
T_{t_i}	0.65	0.4	0.3	0.6	0.8
T_{g_i}	0.1	0.1	0.1	0.1	0.15

	Area 1	Area 2	Area 3	Area 4	Area 5
$\Delta\theta_i$	$\ x_{[1,1]}\ _\infty \leq 0.1$	$\ x_{[2,1]}\ _\infty \leq 0.1$	$\ x_{[3,1]}\ _\infty \leq 0.1$	$\ x_{[4,1]}\ _\infty \leq 0.1$	$\ x_{[5,1]}\ _\infty \leq 0.1$
ΔP_{ref_i}	$\ u_{[1]}\ _\infty \leq 0.5$	$\ u_{[2]}\ _\infty \leq 0.65$	$\ u_{[3]}\ _\infty \leq 0.65$	$\ u_{[4]}\ _\infty \leq 0.55$	$\ u_{[5]}\ _\infty \leq 0.5$

$$P_{12} = 4 \quad P_{23} = 2 \quad P_{34} = 2 \quad P_{45} = 3 \quad P_{25} = 3$$

Table 2: Model parameters and constraints for systems $\Sigma_{[i]}$, $i \in 1 : 5$.

Step time	Area i	ΔP_{L_i}
5	1	+0.15
15	2	-0.15
20	3	+0.12
40	3	-0.12
40	4	+0.28

Table 3: Load of power ΔP_{L_i} (p.u.) for simulation in Scenario 1. $+\Delta P_{L_i}$ means a step of required power, hence a decrease of the frequency deviation $\Delta\omega_i$ and then an increase of the power reference ΔP_{ref_i} , while $-\Delta P_{L_i}$ means the opposite.

Step time	Area i	ΔP_{L_i}
5	1	+0.10
15	2	-0.17
20	1	+0.05
20	2	+0.12
20	3	-0.10
30	3	+0.10
40	4	+0.08
40	5	-0.15

Table 4: Load of power ΔP_{L_i} (p.u.) for simulation in Scenario 2. $+\Delta P_{L_i}$ means a step of required power, hence a decrease of the frequency deviation $\Delta\omega_i$ and then an increase of the power reference ΔP_{ref_i} , while $-\Delta P_{L_i}$ means the opposite.

Step time	Area i	ΔP_{L_i}
5	1	+0.12
15	2	-0.15
20	5	+0.20
40	2	+0.15
40	3	+0.13
40	5	-0.20

Table 5: Load of power ΔP_{L_i} (p.u.) for simulation in Scenario 3. $+\Delta P_{L_i}$ means a step of required power, hence a decrease of the frequency deviation $\Delta\omega_i$ and then an increase of the power reference ΔP_{ref_i} , while $-\Delta P_{L_i}$ means the opposite.