

# Linear programming: sensitivity analysis

G. Ferrari Trecate

Dipartimento di Ingegneria Industriale e dell'Informazione  
Università degli Studi di Pavia

Industrial Automation

# Introduction

Sensitivity analysis is also known as "post-optimality" analysis because it has to be performed **after** an optimal solution has been computed

How the optimizers change when

P 1) the cost changes?

P 2) the vector  $b$  changes?

Entries in the matrix  $A$  could change as well but this will not be covered in this class.

## Ex. Product mix

$$\max x \quad 30x_1 + 20x_2$$

$$8x_1 + 4x_2 \leq 640$$

$$4x_1 + 6x_2 \leq 540$$

$$x_1 + x_2 \leq 100$$

$$x_1, x_2 \geq 0$$

$$x^* = \begin{bmatrix} x_1^* \\ x_2^* \end{bmatrix} = \begin{bmatrix} 60 \\ 40 \end{bmatrix}$$

(P1)  $\longleftrightarrow$  robustness of the optimal production plan when revenues are uncertain

(P2)  $\longleftrightarrow$  is it convenient to buy additional resources (i.e. man/hours in the production phases)?

Ex. of problem P1

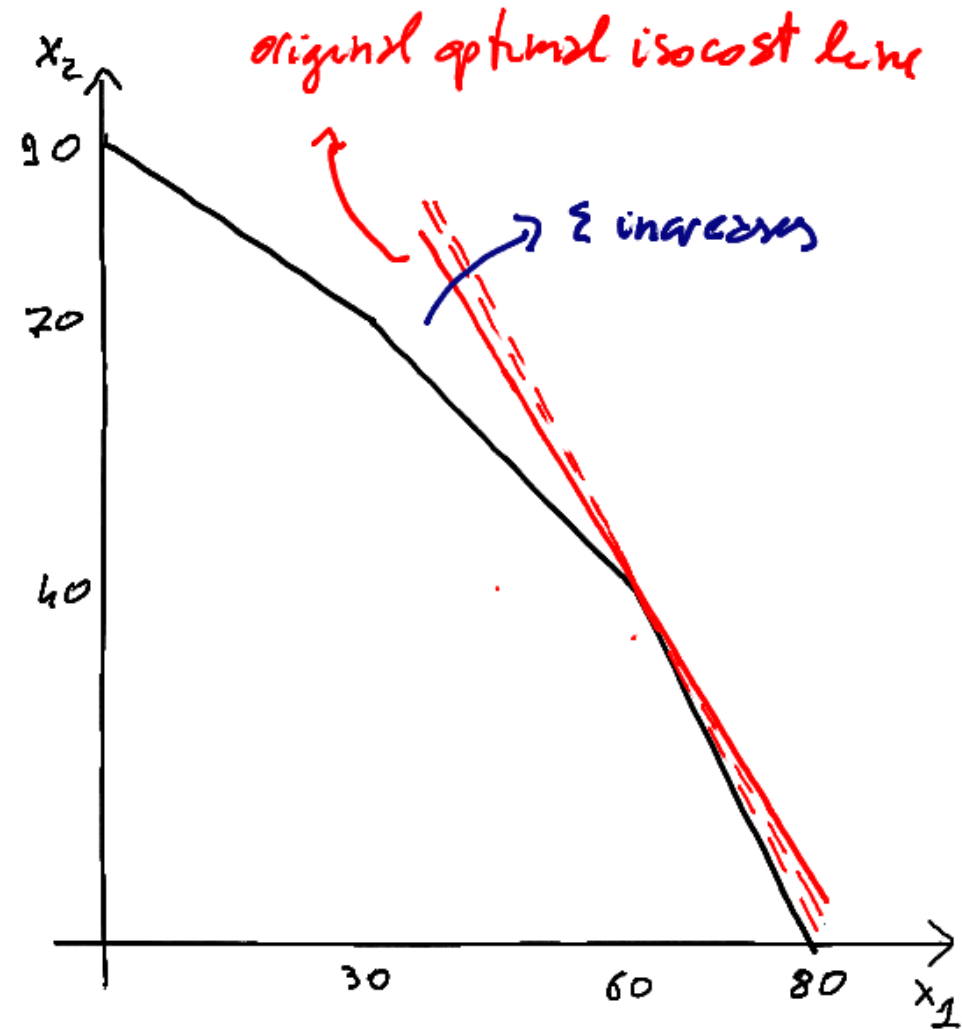
$$\bar{c} = [c_1 \ c_2] = [30 \ 20]$$

How much  $c_1$  can be increased without changing the optimal solution?

$$c_{\perp}: (30 + \varepsilon) x_1 + 20 x_2 = \alpha$$

$$\hookrightarrow x_2 = \frac{\alpha}{20} - \frac{30 + \varepsilon}{20} x_1$$

$\varepsilon > 0$  increases  $\rightarrow c_{\perp}$  rotates clockwise  
 $\rightarrow$  the optimal solution does not change until the optimal  $c_{\perp}$  intersects also the vertex  $[80, 0]^T$



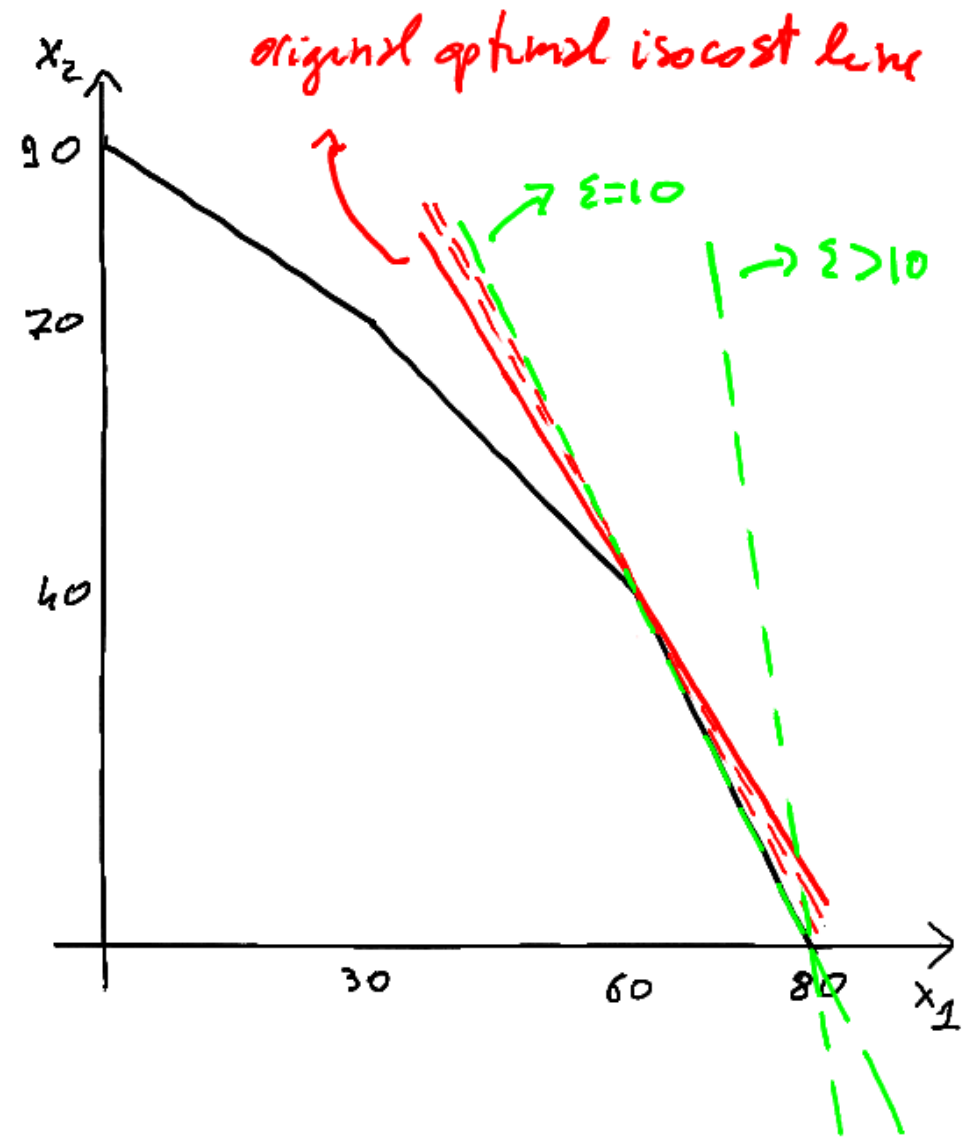
Compute the maximal  $\varepsilon$

$$x^T = [60, 40] \rightarrow (30 + \varepsilon)60 + 20 \cdot 40 = \alpha$$

$$x^T = [80, 0] \rightarrow (30 + \varepsilon)80 = \alpha$$

$$\varepsilon = 10 \text{ and } \alpha = 3200$$

- $\varepsilon = 10$ : the optimal solution is not unique
- $\varepsilon > 10$ : the optimal solution becomes  $[80, 0]^T$



How much  $c_2$  can be decreased without changing the optimal solution?

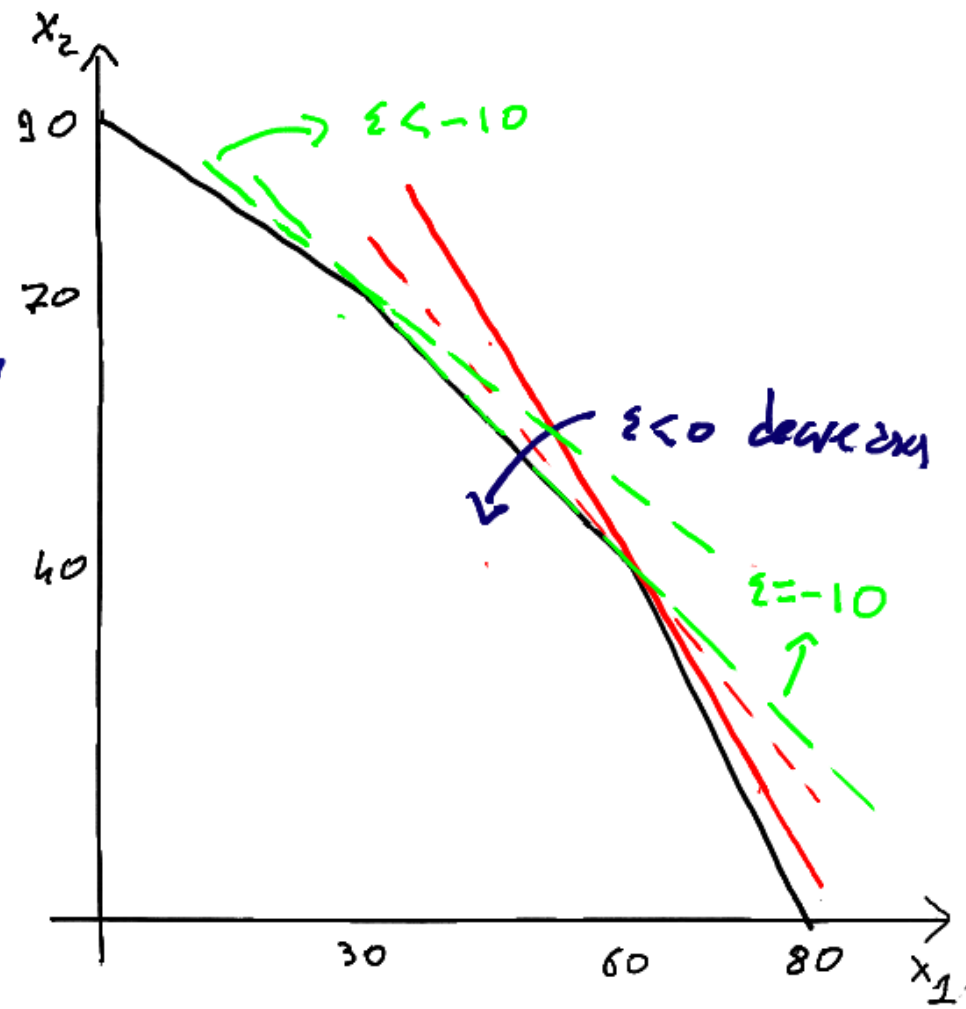
$\epsilon < 0$  decreases  $\rightarrow c_2$  rotates counter clockwise  $\rightarrow$  the optimal solution does not change until  $c_2$  intersects  $[30 \ 70]^T$ .

$\therefore$  same computations as before

This happens for  $\epsilon = -10$

**Conclusion:** for  $c_1 \in [20, 40]$  the optimal solution does not change

Ex. @ home: same analysis for  $c_1 = 30$  and  $c_2 = 20 + \epsilon$



## Ex. of problem P2

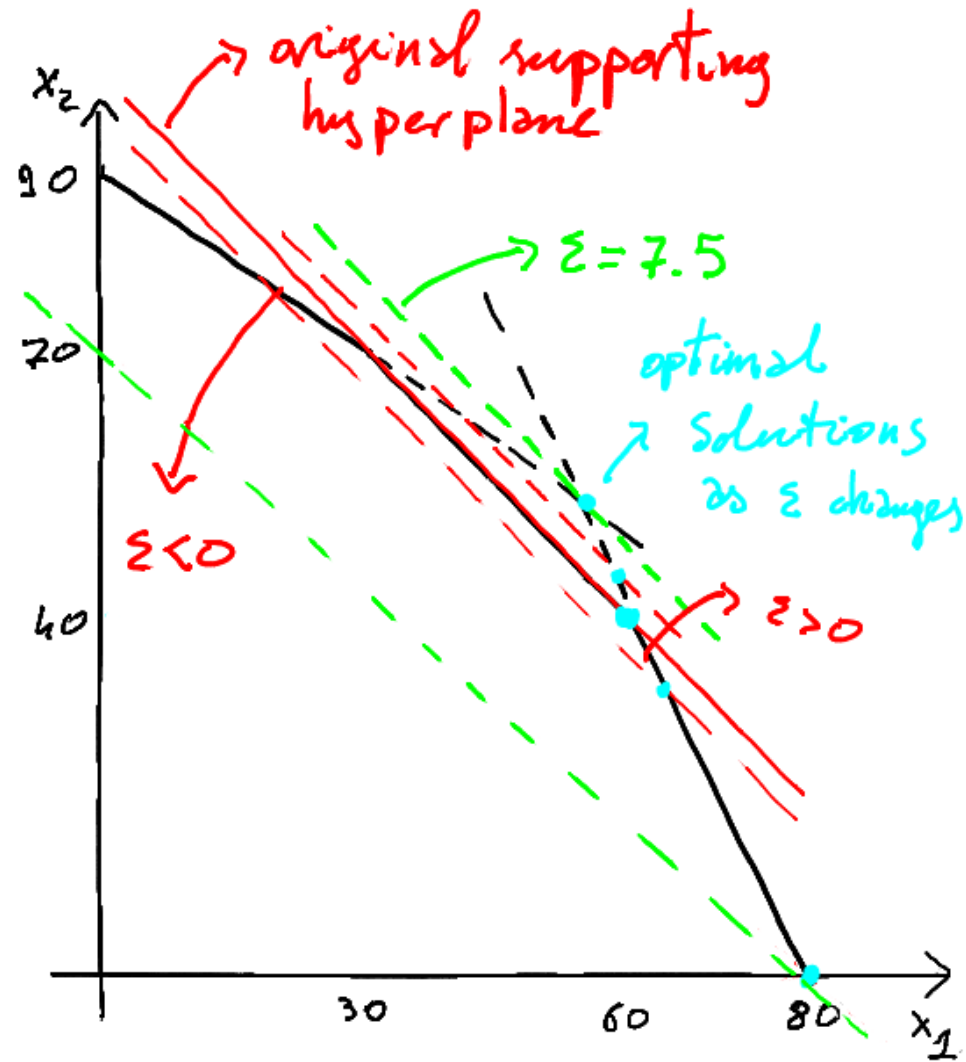
Change of resources in the quality control phase: the constraint  $x_1 + x_2 \leq 100$  gets replaced by

$$x_1 + x_2 \leq 100 + \varepsilon$$

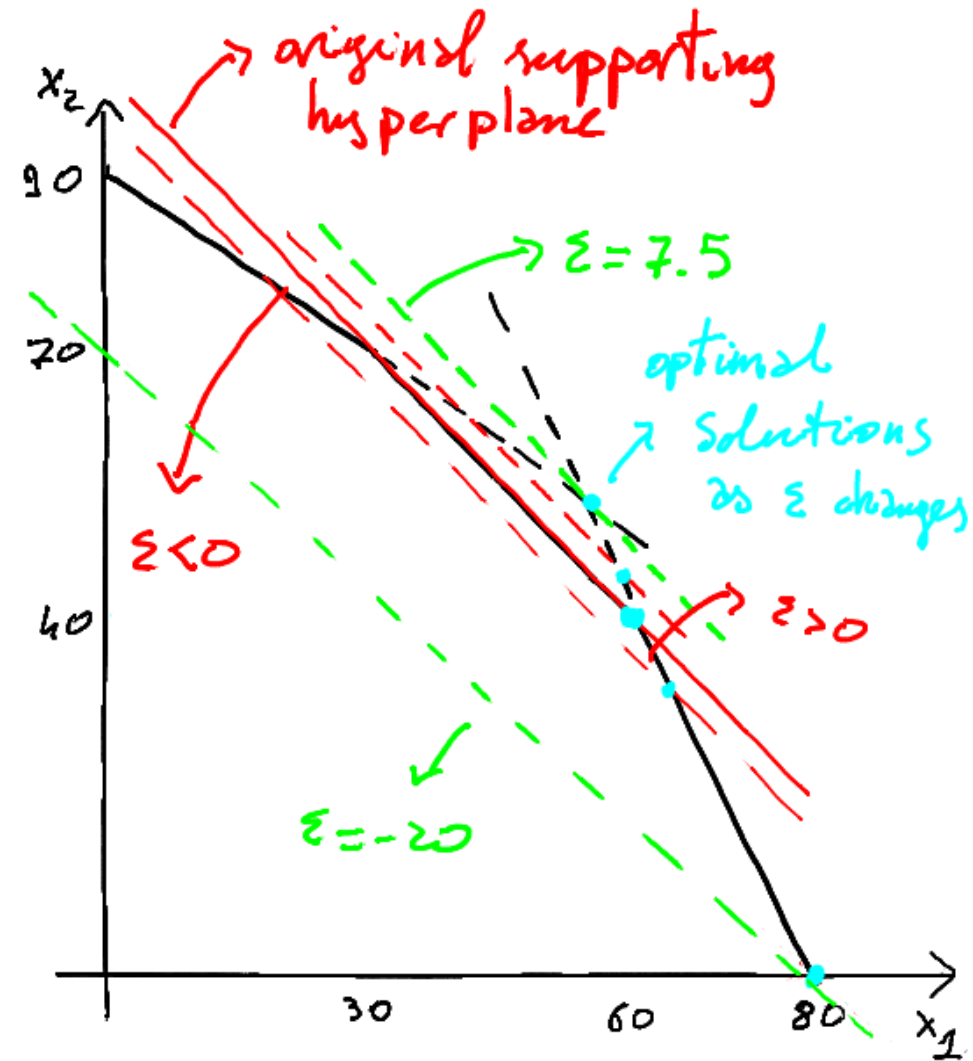
Supporting hyperplane:  $x_2 = (100 + \varepsilon) - x_1$

The optimal solution changes for  $\varepsilon \neq 0$  but it lies on the same supporting hyperplanes if  $\varepsilon$  does not vary too much

↳ The optimal basis does not change



- $\varepsilon < 0$ : the optimal basis does not change until  $\varepsilon \in [-20, 0]$   
 $\hookrightarrow \varepsilon = -20$ : the supporting hyperplanes intersect  $[80, 0]^T$
- $\varepsilon > 0$ : the optimal basis does not change until  $\varepsilon \in [0, 7.5]$   
 $\hookrightarrow \varepsilon = 7.5$ : the constraint becomes redundant





# Sensitivity analysis: general theory

Reference problem: standard LP

$$\begin{aligned} z^* &= \max c^T x \\ Ax &= b \\ x &\geq 0 \end{aligned}$$

Optimal BFS given by  $x_B$  and  $x_F = 0$ . Assumption: it is not degenerate.

Optimality conditions

$$(01): x_B = B^{-1}b \geq 0 \quad (\text{feasibility})$$

$$(02): r_F^T = c_F^T - c_B^T B^{-1}F \leq 0 \quad (\text{optimality})$$

Sensitivity analysis: study when B is still optimal as

- (a) only  $b$  changes
- (b) only  $c_B$  changes
- (c) only  $c_F$  changes

Case (a):  $b \rightarrow b + \Delta b$

$$(O1): x_B = B^{-1}(b + \Delta b) \geq 0$$

(O2) does not change

$B$  does not change until  $\Delta b$  verifies

$$B^{-1}b \geq -B^{-1}\Delta b$$

→ admissible perturbations  
lie in a polyhedron

However the optimal solution changes and the optimal cost becomes

$$c_B^T B^{-1}(b + \Delta b) = (\lambda^*)^T (b + \Delta b) \quad \lambda^* = \text{optimal multipliers}$$

**Rmk.**  $\lambda_i^*$  measures the sensitivity of the optimal cost to  $\Delta b_i$

Case (b):  $c_B \rightarrow c_B + \Delta c_B$

Let  $\bar{r}_F^T = c_F^T - c_B^T B^{-1} F$  (reduced cost for  $\Delta c_B = 0$ )

(O1) does not change

$$(O2) \quad c_F^T - (c_B^T + \Delta c_B^T) B^{-1} F \leq 0$$

$$\hookrightarrow \underbrace{c_F^T - c_B^T B^{-1} F}_{\bar{r}_F^T} - \Delta c_B^T B^{-1} F \leq 0$$

The optimal basis (and the optimal solution) does not change if  $\Delta c_B$  verifies

$$\Delta c_B^T B^{-1} F \geq \bar{r}_F^T$$

$\hookrightarrow$  admissible perturbations lie in a polyhedron

Case (c):  $c_F \rightarrow c_F + \Delta c_F$

(O1) does not change

$$(O2): c_F^T + \Delta c_F^T - c_B^T B^{-1} F \leq 0$$

The optimal basis (and the optimal solution) does not change if  $\Delta c_F$  verifies

$$\Delta c_F \leq -\bar{r}_F$$

**Rmk.** For "min" problems the inequalities in cases (b) and (c) must be reversed

Ex.

$$\max [3 \ 1 \ 3 \ 0 \ 0 \ 0] x$$

$$x = [x_1 \ x_2 \ x_3 \ x_4 \ x_5 \ x_6]^T$$

$$\begin{bmatrix} 2 & 1 & 1 & 1 & 0 & 0 \\ 1 & 2 & 3 & 0 & 1 & 0 \\ 2 & 2 & 1 & 0 & 0 & 1 \end{bmatrix} x = \begin{bmatrix} 2 \\ 5 \\ 6 \end{bmatrix}$$

$$x \geq 0$$

The simplex algorithm produces the optimal tableau

		$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$
	$-\frac{22}{5}$	0	$\frac{7}{5}$	0	$-\frac{6}{5}$	$-\frac{3}{5}$	0
$x_1$	$\frac{1}{5}$	1	$\frac{1}{5}$	0	$\frac{3}{5}$	$-\frac{1}{5}$	0
$x_3$	$\frac{8}{5}$	0	$\frac{3}{5}$	1	$-\frac{1}{5}$	$\frac{2}{5}$	0
$x_6$	4	0	1	0	-1	0	1

Is B still optimal when

(i)  $b_1 \rightarrow b_1 + 5$

(ii)  $c_3 \rightarrow c_3 + 1$  ?

Useful quantities from the optimal tableau

		$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$
	$-\frac{27}{5}$	0	$-\frac{7}{5}$	0	$-\frac{6}{5}$	$-\frac{3}{5}$	0
$x_1$	$\frac{1}{5}$	1	$\frac{1}{5}$	0	$\frac{3}{5}$	$-\frac{1}{5}$	0
$x_3$	$\frac{8}{5}$	0	$\frac{3}{5}$	1	$-\frac{1}{5}$	$\frac{2}{5}$	0
$x_6$	4	0	1	0	-1	0	1

$$x_B = \begin{bmatrix} x_1 \\ x_3 \\ x_6 \end{bmatrix}$$

$$B = \begin{bmatrix} 2 & 1 & 0 \\ 1 & 3 & 0 \\ 2 & 1 & 1 \end{bmatrix}$$

$$x_B^* = B^{-1}b = \begin{bmatrix} 1/5 \\ 8/5 \\ 4 \end{bmatrix}$$

$$x_F = \begin{bmatrix} x_2 \\ x_4 \\ x_5 \end{bmatrix}$$

$$\bar{r}_F = \begin{bmatrix} -7/5 \\ -6/5 \\ -3/5 \end{bmatrix}$$

$$B^{-1} = \begin{bmatrix} 3/5 & -1/5 & 0 \\ -1/5 & 2/5 & 0 \\ -1 & 0 & 1 \end{bmatrix}$$

$$\bar{F} = B^{-1}F = \begin{bmatrix} 1/5 & 3/5 & -2/5 \\ 3/5 & -1/5 & 2/5 \\ 1 & -1 & 0 \end{bmatrix}$$

Problem (i):  $b_2 \rightarrow b_2 + 5$

$B$  optimal  $\Leftrightarrow B^{-1}b \geq -B^{-1}\Delta b$

$$\Delta b = \begin{bmatrix} \Delta b_1 \\ 0 \\ 0 \end{bmatrix} \quad \begin{bmatrix} +\frac{1}{5} \\ \frac{8}{5} \\ 4 \end{bmatrix} \geq \begin{bmatrix} -\frac{3}{5} \Delta b_1 \\ \frac{1}{5} \Delta b_1 \\ \Delta b_1 \end{bmatrix} \quad \rightarrow \quad \begin{cases} \Delta b_1 \geq -\frac{1}{3} \\ \Delta b_1 \leq 8 \\ \Delta b_1 \leq 4 \end{cases}$$

$B$  does not change if  $\Delta b_1 \in [-\frac{1}{3}, 4]$ .

For  $\Delta b_1 = 5$ ,  $B$  is no longer optimal

@ home: is  $B$  still optimal when  $b_2 \rightarrow b_2 - 1$  or  $b_3 \rightarrow b_3 - 1$ ?

Problem (ii):  $c_3 \rightarrow c_3 + 1$

$c_3$  is an element of  $c_B$

$B$  is still optimal only if  $\Delta c_B^T B^{-1} F \geq \bar{r}_F^T$

$$[\Delta c_1 \quad \Delta c_3 \quad \Delta c_6] \bar{F} \geq \left[ -\frac{7}{5} \quad -\frac{6}{5} \quad -\frac{3}{5} \right]$$

$$\left. \begin{array}{l} \frac{1}{5} \Delta c_1 + \frac{3}{5} \Delta c_3 + \Delta c_6 \geq -\frac{7}{5} \\ \frac{3}{5} \Delta c_1 - \frac{1}{5} \Delta c_3 - \Delta c_6 \geq -\frac{6}{5} \\ -\frac{1}{5} \Delta c_1 + \frac{2}{5} \Delta c_3 \geq -\frac{3}{5} \end{array} \right\} \xrightarrow{\Delta c_1 = \Delta c_6 = 0} \left. \begin{array}{l} \Delta c_3 \geq -\frac{7}{5} \\ \Delta c_3 \leq 6 \\ \Delta c_3 \geq -\frac{3}{2} \end{array} \right\}$$

$B$  does not change  
if  $\Delta c_3 \in \left[ -\frac{3}{2}, 6 \right]$   
 $\downarrow$   
 $B$  is still optimal  
for  $c_3 \rightarrow c_3 + 1$



Next: exercises!